

Quantum...

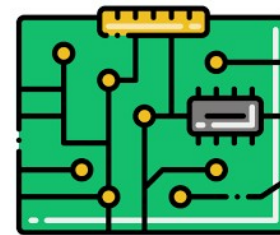
Computing



God,



Gates and Circuits



HOW'S YOUR
QUANTUM COMPUTER
PROTOTYPE COMING
ALONG?

GREAT!

THE PROJECT EXISTS
IN A SIMULTANEOUS
STATE OF BEING BOTH
TOTALLY SUCCESSFUL
AND NOT EVEN
STARTED.

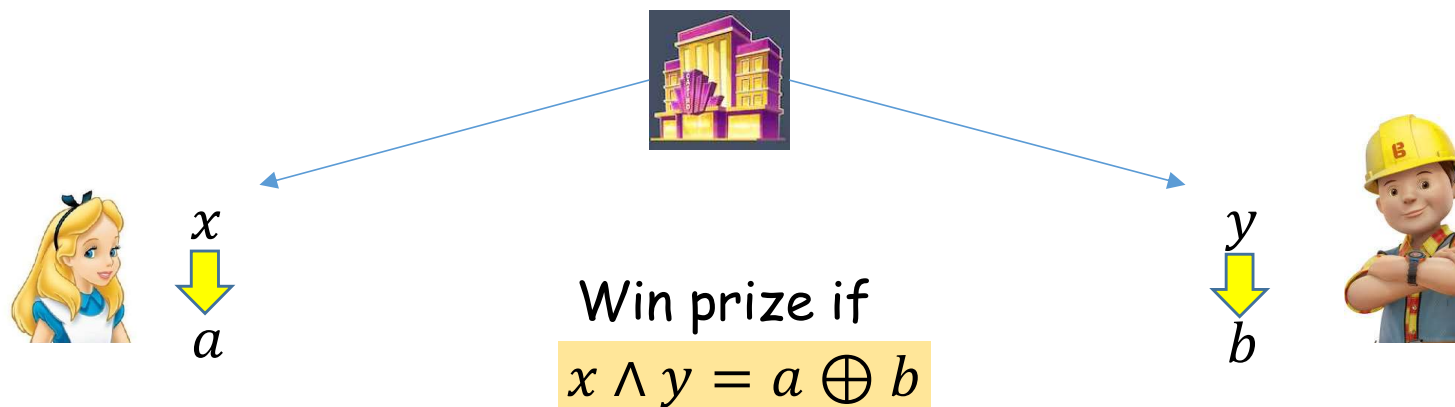
CAN I
OBSERVE
IT?

THAT'S
A TRICKY
QUESTION.

Dilbert.com DilbertCartoonist@gmail.com

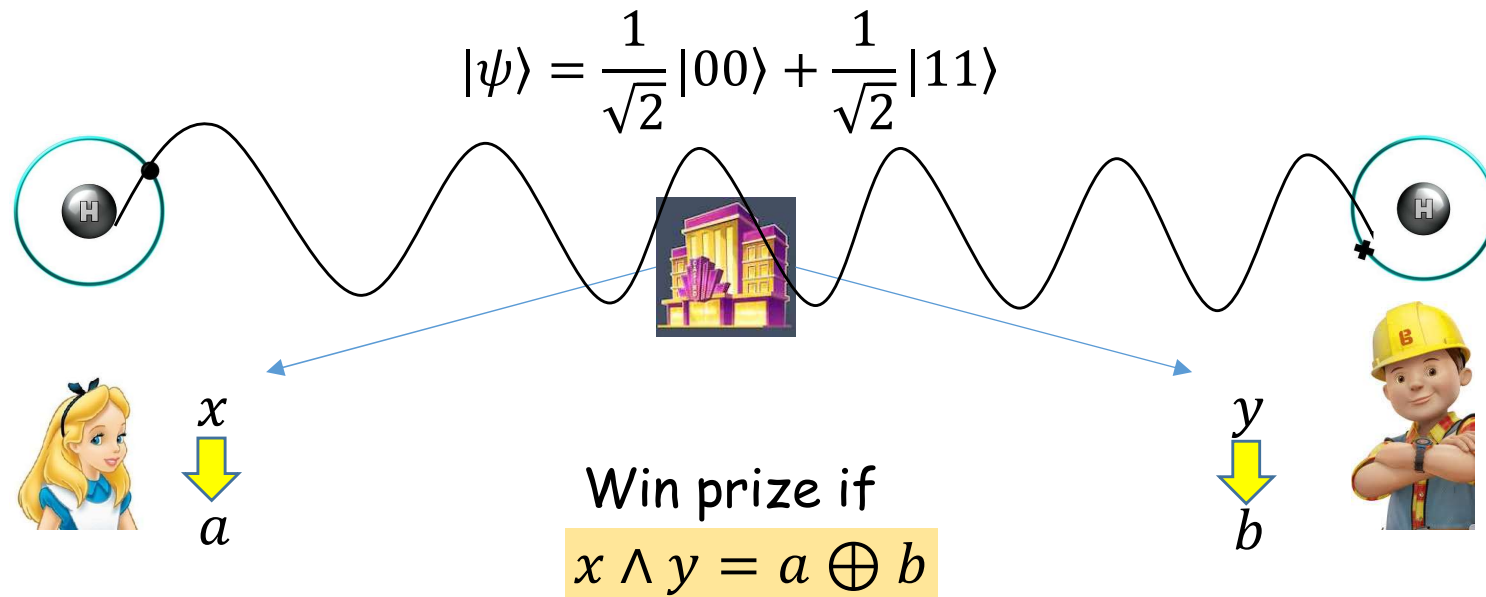
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The CHSH game



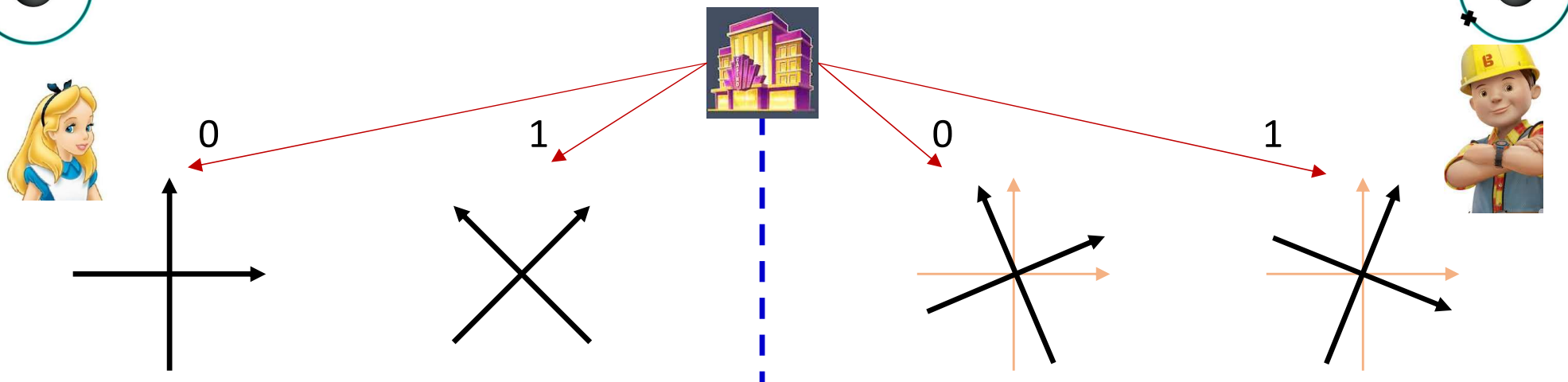
- Alice and Bob live in separate cities and may not communicate
- The casino sends each of them a random bit
 - Need not be identical
- They must inspect their bit and output a value
 - Alice outputs a , Bob outputs b
- They get a prize of \$1.00 if:
 - Both got “1” from the casino and their outputs are such that $a \neq b$
 - Any other condition ($[0,1]$, $[1,0]$, $[0,0]$) they must output $a == b$
- What is the best strategy, and what is their expected earning?

The CHSH game with a qubit



- Before moving to separate cities, Alice and Bob split a pair of entangled bits in the Bell State
- Now what is their best strategy?

The CHSH game with a qubit



- Alice uses two sets of bases for measurement: 0/1 and +/- (at 45°)
 - If Alice gets a 0 from the casino she measures using 0/1 and outputs the value
 - Else she measures using +/- and outputs the value
- Bob uses two sets of bases: at $\left[\frac{\pi}{8}, \frac{5\pi}{8}\right]$ and $\left[\frac{-\pi}{8}, \frac{3\pi}{8}\right]$
 - If Bob gets a 0 from the casino, he measures using $\left[\frac{\pi}{8}, \frac{5\pi}{8}\right]$ and outputs the value
 - Else he measures using the $\left[\frac{-\pi}{8}, \frac{3\pi}{8}\right]$ and outputs the value

The CHSH inequality

- The Clauser Horne Shimony Holt (1969):
- For classical computers

$$E[0,0] + E[0,1] + E[1,0] - E[1,1] \leq 2$$

- where $E[x, y]$ is the probability that Alice and Bob “agree” (i.e. $a = b$) when they receive x and y respectively
 - Note: The maximum possible value under perfect knowledge is 3. The closer you are to 3, the more money you make

- Using quantum entanglement

$$E[0,0] + E[0,1] + E[1,0] - E[1,1] \leq 2\sqrt{2}$$

- Regardless of the actual qubit shared
 - Over any policy / measurement strategy
 - This is 2.8, which is very close to the max possible value of 3
- Qubits, which are useless for communication, can still be used to create *correlations* which can be exploited
 - They can “enhance” asymmetries in the system

Lesson – you cannot communicate

- But you can correlate
- And correlation can be used for profit...

The Determinism Conundrum

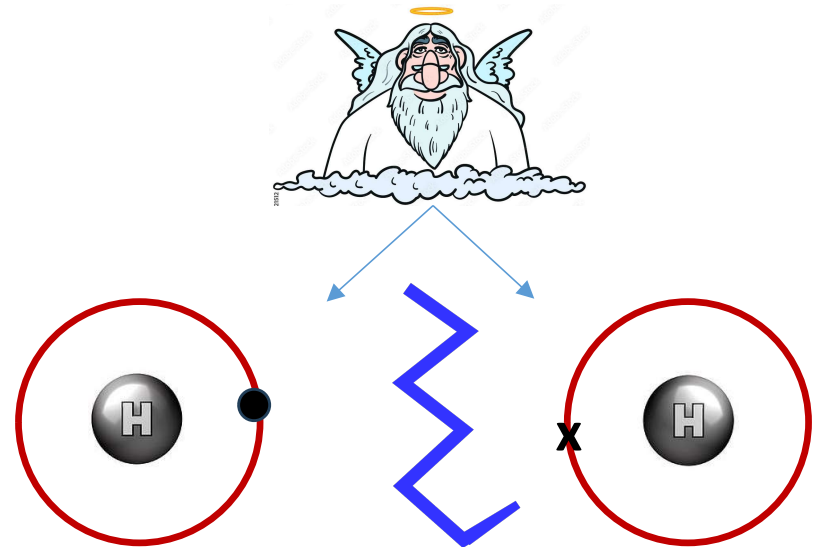
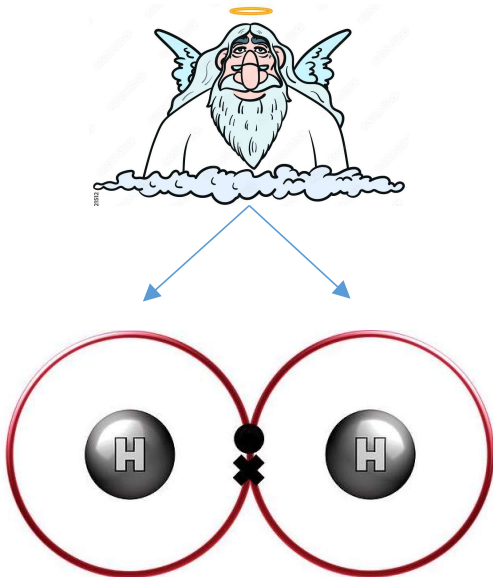
- Schroedingers equation

$$i\hbar \frac{\partial}{\partial t} \Psi(\mathbf{x}, t) = \left[-\frac{\hbar^2}{2m} \nabla^2 + V(\mathbf{x}, t) \right] \Psi(\mathbf{x}, t)$$

- The term in the square brackets is the *Hamiltonian*.
 - It includes a *potential* term $V(\mathbf{x}, t)$ which can be manipulated to manipulate the Hamiltonian itself
- ∇^2 is the *Laplacian*
- $\mathbf{x}$ is 3D position
- $\Psi(\mathbf{x}, t)$ is the wave function for a particle

Is there a Higher Knower?

- $\Psi(x, t)$, when measured, collapses into one of the many possible states
- *Was this state fore-ordained?*
- Einstein, Podolsky and Rosen (EPR) – yes!
 - **Local realism:** The fate of one qubit cannot affect another faster than light
 - Ergo: The “entangled” qubits were *foreordained* to their state by some (possibly ancient) latent variable. There is no “entanglement” per-se...



Is there a Higher Knower?

- $\Psi(x, t)$, when measured, collapses into one of the many possible states
- *Was this state fore-ordained?*
- The CHSH game uses entangled qubits to prove otherwise
 - Einstein, Podolsky, Rosen: The two cubits were independently fore-ordained by a common cause to fall the same way
 - Bell: If so, their individual measurements (and their bases) should not have any influence on the other if they are randomly chosen
 - $P(\text{match}) \leq 0.75$
 - $P = 0.85$ implies actual entanglement, and genuine randomness in measurement

What does the higher-knower know?

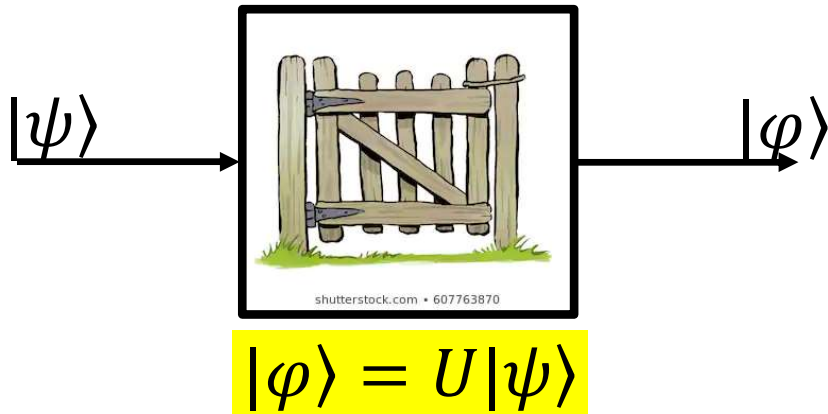
- The *probability distribution* in $\Psi(\mathbf{x}, t)$ may be known for all $\mathbf{x}$ at some t .

- But

$$i\hbar \frac{\partial}{\partial t} \Psi(\mathbf{x}, t) = \left[-\frac{\hbar^2}{2m} \nabla^2 + V(\mathbf{x}, t) \right] \Psi(\mathbf{x}, t)$$

- This is solveable.
- If $\Psi(\mathbf{x}, t)$ is known at any time it is known for all time!!!
 - The wavefunction is fully determined for all time
 - The universe is deterministic in probability

Quantum gates



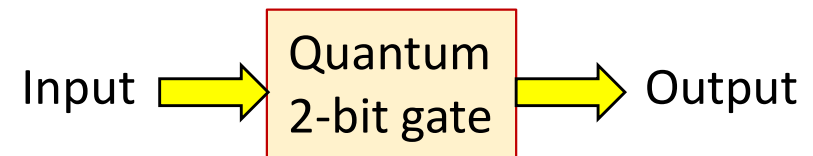
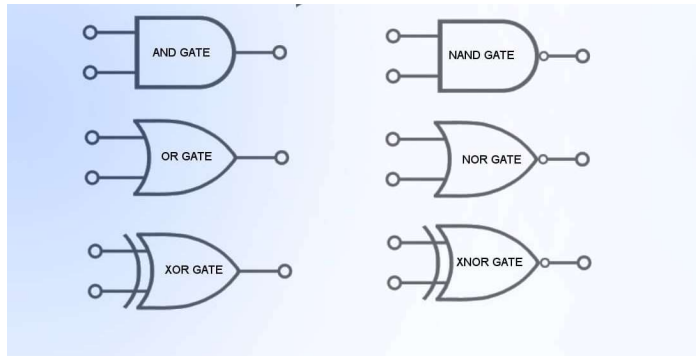
$$|\psi\rangle = \begin{bmatrix} a_0 \\ \vdots \\ a_N \end{bmatrix} \quad |\varphi\rangle = \begin{bmatrix} b_0 \\ \vdots \\ b_N \end{bmatrix}$$

$$\begin{bmatrix} b_0 \\ \vdots \\ b_N \end{bmatrix} = U \begin{bmatrix} a_0 \\ \vdots \\ a_N \end{bmatrix}$$

- Operate on N -Dimensional phasors
 - In a complex N -D complex Hilbert space
 - The input is an N -D phasor, the output too is an N -D phasor
- The “gate” is itself a transform
 - **A unitary transform**
- So how many inputs does the gate have, and how many outputs?

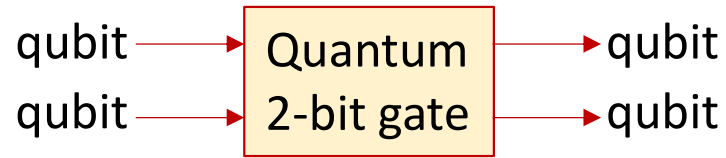
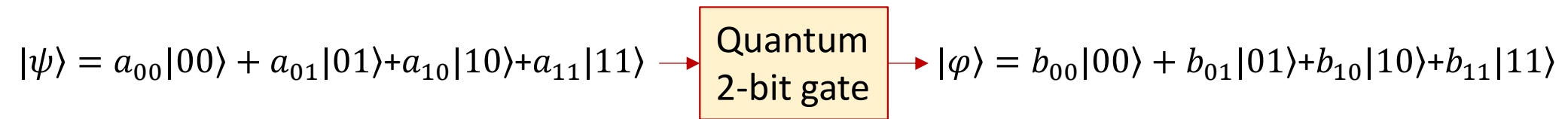
TWO QUBIT GATES

Classical two-bit gates



- Two-bit classical gates take in 2 bits, and output one or two bits
 - They operate on a two-dimensional input space
- What dimensionality of space do two-qubit quantum gates operate on

TWO QUBIT GATES



- Two-bit classical gates take in 2 bits, and output one or two bits
 - They operate on a two-dimensional input space
- What dimensionality of space do two-qubit quantum gates operate on
 - Four dimensional inputs and outputs
 - But physically still represented by two qubits (thanks quantum!)

Two-qubit gates

$$|\psi\rangle = a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle \rightarrow \boxed{\text{Quantum 2-bit gate}} \rightarrow |\varphi\rangle = b_{00}|00\rangle + b_{01}|01\rangle + b_{10}|10\rangle + b_{11}|11\rangle$$

U

$$\begin{bmatrix} b_{00} \\ b_{01} \\ b_{10} \\ b_{11} \end{bmatrix} = \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ u_{10} & u_{11} & u_{12} & u_{13} \\ u_{20} & u_{21} & u_{22} & u_{23} \\ u_{30} & u_{31} & u_{32} & u_{33} \end{bmatrix} \begin{bmatrix} a_{00} \\ a_{01} \\ a_{10} \\ a_{11} \end{bmatrix}$$

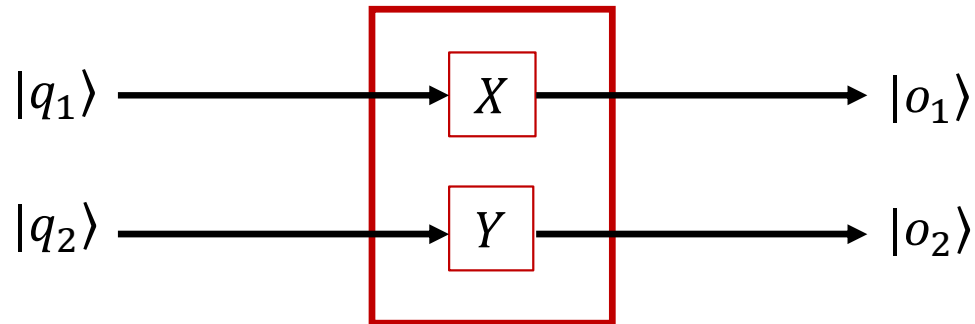
- A single two-qubit gate U operates on the phasor

$$a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle \text{ to output } b_{00}|00\rangle + b_{01}|01\rangle + b_{10}|10\rangle + b_{11}|11\rangle$$

$$U(a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle) = b_{00}|00\rangle + b_{01}|01\rangle + b_{10}|10\rangle + b_{11}|11\rangle$$

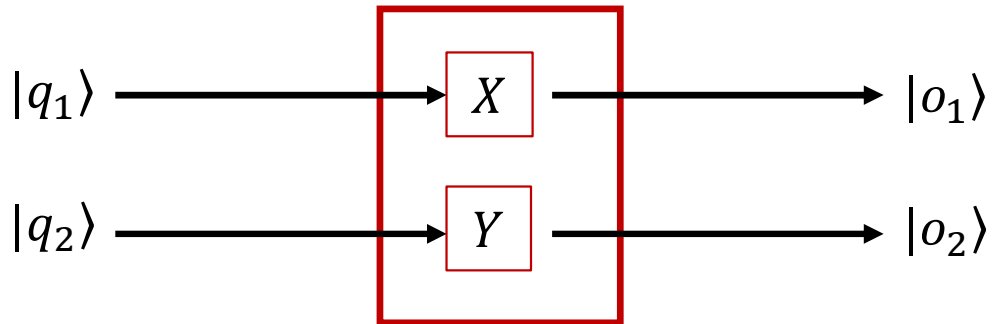
- U is a 4x4 matrix.
 - It is a unitary transform
 - Its columns are orthogonal and form a new basis set
- Examples coming up

The simplest 2-qubit gate



- Each qubit is independently operated on by a one-qubit gate
 - Note: This is still a two qubit system operating on two qubits and producing two qubits
- What is the resulting 2-qubit gate?

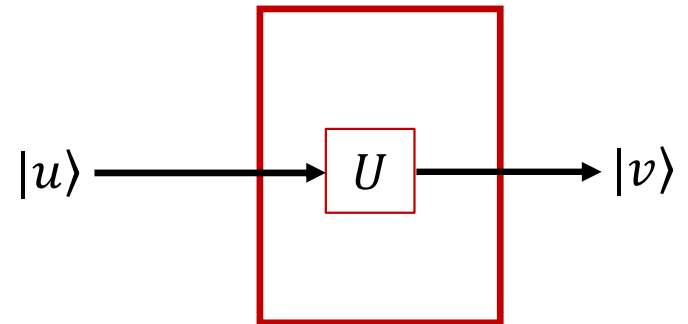
2-qubit gate



$$\begin{aligned} |q_1\rangle &= a_0|0\rangle + a_1|1\rangle & |o_1\rangle &= c_0|0\rangle + c_1|1\rangle \\ |q_2\rangle &= b_0|0\rangle + b_1|1\rangle & |o_2\rangle &= d_0|0\rangle + d_1|1\rangle \end{aligned}$$

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} x_{00} & x_{01} \\ x_{10} & x_{11} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \end{bmatrix}$$

$$\begin{bmatrix} d_0 \\ d_1 \end{bmatrix} = \begin{bmatrix} y_{00} & y_{01} \\ y_{10} & y_{11} \end{bmatrix} \begin{bmatrix} b_0 \\ b_1 \end{bmatrix}$$



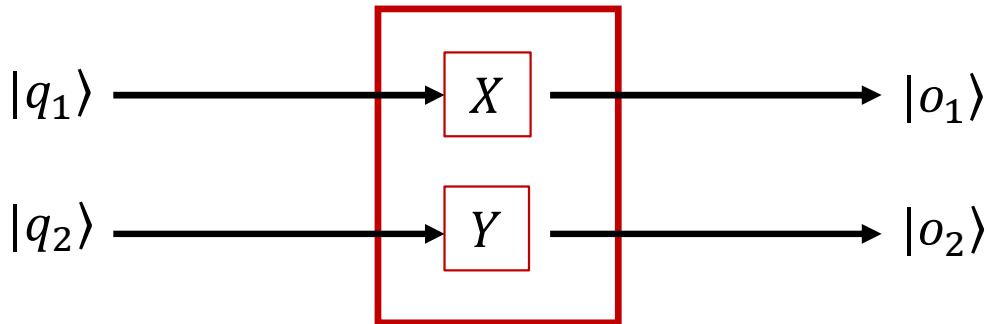
$$|u\rangle = |q_1\rangle \otimes |q_2\rangle = |q_1 q_2\rangle$$

$$|v\rangle = |o_1\rangle \otimes |o_2\rangle = |o_1 o_2\rangle$$

$$\begin{bmatrix} v_{00} \\ v_{01} \\ v_{10} \\ v_{11} \end{bmatrix} = \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ u_{10} & u_{11} & u_{12} & u_{13} \\ u_{20} & u_{21} & u_{22} & u_{23} \\ u_{30} & u_{31} & u_{32} & u_{33} \end{bmatrix} \begin{bmatrix} u_{00} \\ u_{01} \\ u_{10} \\ u_{11} \end{bmatrix}$$

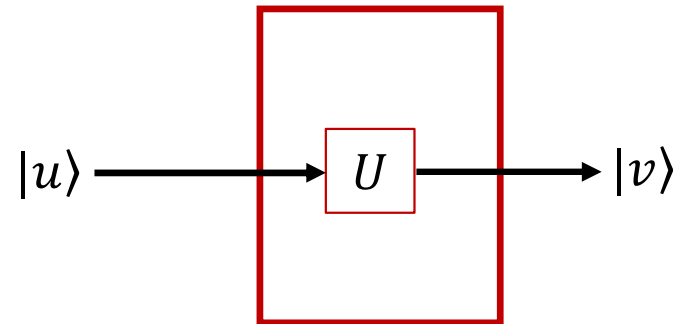
- How does U relate to X and Y

2-qubit gate



$$X = \begin{bmatrix} x_{00} & x_{01} \\ x_{10} & x_{11} \end{bmatrix}$$

$$Y = \begin{bmatrix} y_{00} & y_{01} \\ y_{10} & y_{11} \end{bmatrix}$$



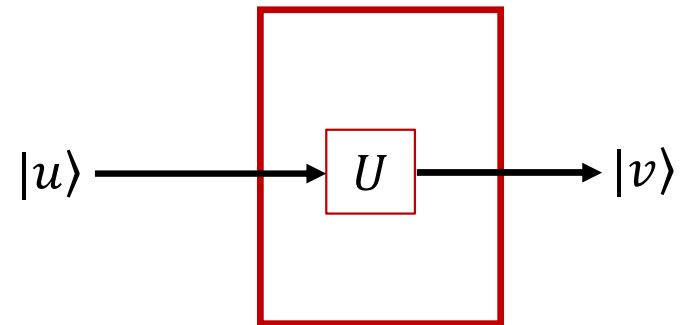
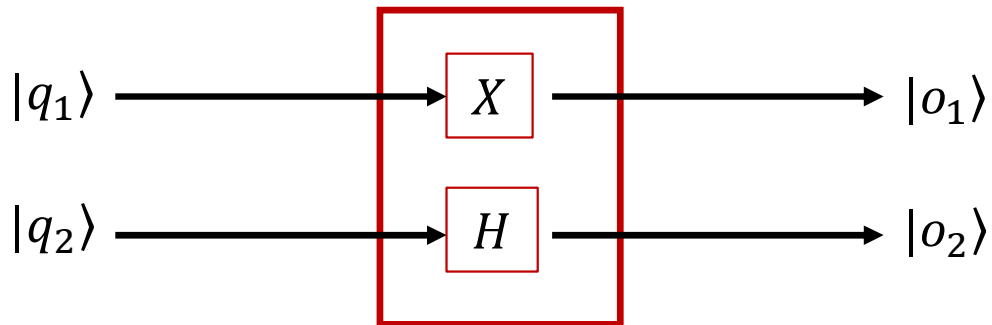
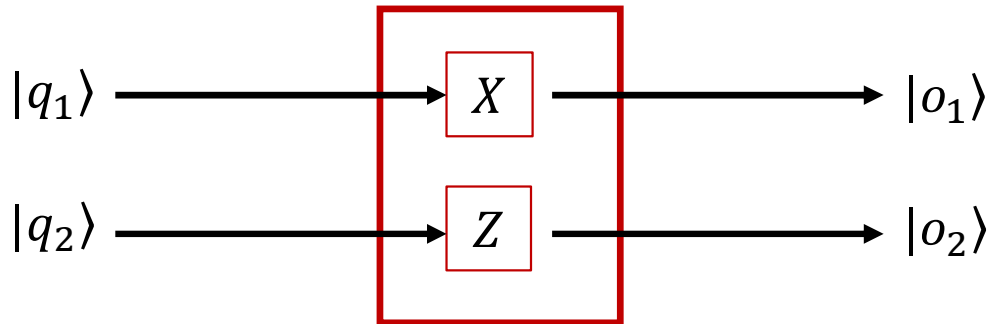
$$U = X \otimes Y$$

$$U = X \otimes Y = \begin{bmatrix} x_{00} \begin{bmatrix} y_{00} & y_{01} \\ y_{10} & y_{11} \end{bmatrix} & x_{01} \begin{bmatrix} y_{00} & y_{01} \\ y_{10} & y_{11} \end{bmatrix} \\ x_{10} \begin{bmatrix} y_{00} & y_{01} \\ y_{10} & y_{11} \end{bmatrix} & x_{11} \begin{bmatrix} y_{00} & y_{01} \\ y_{10} & y_{11} \end{bmatrix} \end{bmatrix}$$

- How does U relate to X and Y when the qubits don't interact
 - I.e. no entanglement
 - Verify that U is unitary

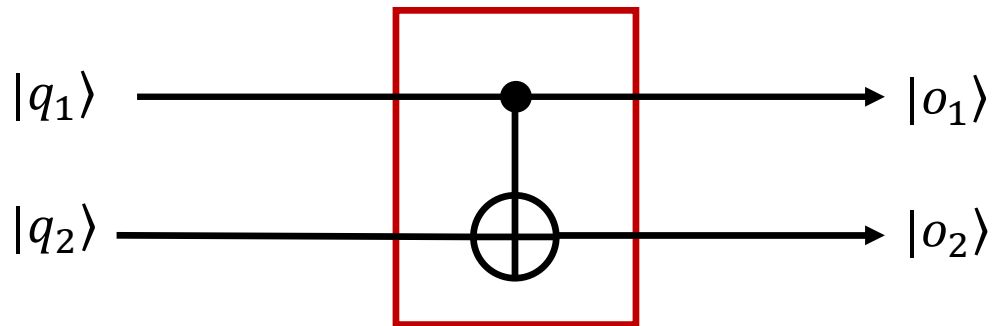
$$U = \begin{bmatrix} x_{00}y_{00} & x_{00}y_{01} & x_{01}y_{00} & x_{01}y_{01} \\ x_{00}y_{10} & x_{00}y_{11} & x_{01}y_{10} & x_{01}y_{11} \\ x_{10}y_{00} & x_{10}y_{01} & x_{11}y_{00} & x_{11}y_{01} \\ x_{10}y_{10} & x_{10}y_{11} & x_{11}y_{10} & x_{11}y_{11} \end{bmatrix}$$

2-qubit gate



- What is U ?
- What is the structure of the transform?
 - Can you generalize to more than 2 non-interacting bits?

The CNOT gate

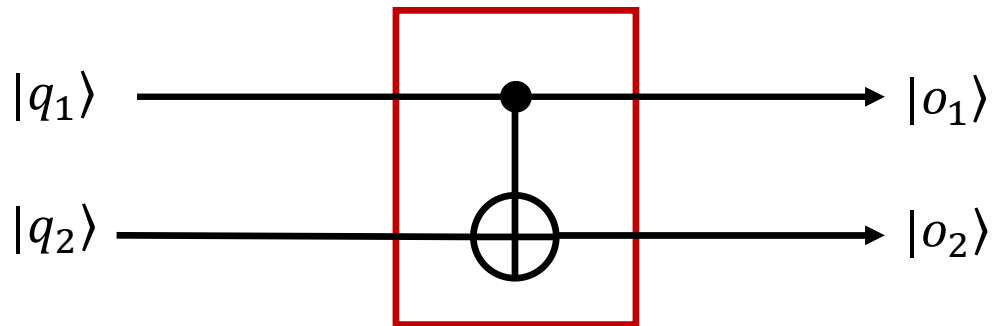


$$|o_1 o_2\rangle = CNOT(|q_1 q_2\rangle)$$

$$o_1 = q_1 \quad o_2 = q_1 \oplus q_2$$

- $o_1 = q_1$
 - The first input is always unchanged
- $o_2 = q_2$ if $q_1 = |0\rangle$ (is 0), otherwise $o_2 = X(q_2)$
 - It bit-flips if the first input is 1
- **q_1 is the *control* bit and q_2 is the *target* bit**

The CNOT gate



$$|o_1 o_2\rangle = CNOT(|q_1 q_2\rangle)$$

$$\begin{bmatrix} u_{00} \\ u_{01} \\ u_{11} \\ u_{10} \end{bmatrix} = [CNOT] \begin{bmatrix} u_{00} \\ u_{01} \\ u_{10} \\ u_{11} \end{bmatrix}$$

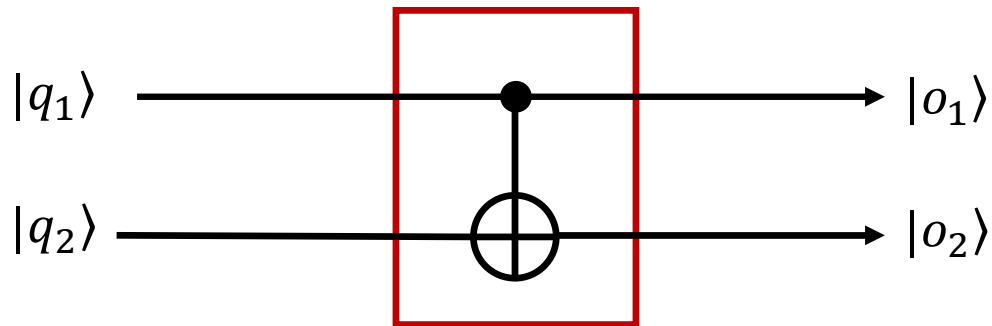
$$CNOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

Verify that this is unitary

$$CNOT(a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle) = a_{00}|00\rangle + a_{01}|01\rangle + a_{11}|10\rangle + a_{10}|11\rangle$$

- Are o_1 and o_2 entangled?

The CNOT gate



$$|o_1 o_2\rangle = CNOT(|q_1 q_2\rangle)$$

$$\begin{bmatrix} u_{00} \\ u_{01} \\ u_{11} \\ u_{10} \end{bmatrix} = [CNOT] \begin{bmatrix} u_{00} \\ u_{01} \\ u_{10} \\ u_{11} \end{bmatrix}$$

$$CNOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$CNOT(a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle) = a_{00}|00\rangle + a_{01}|01\rangle + a_{11}|10\rangle + a_{10}|11\rangle$$

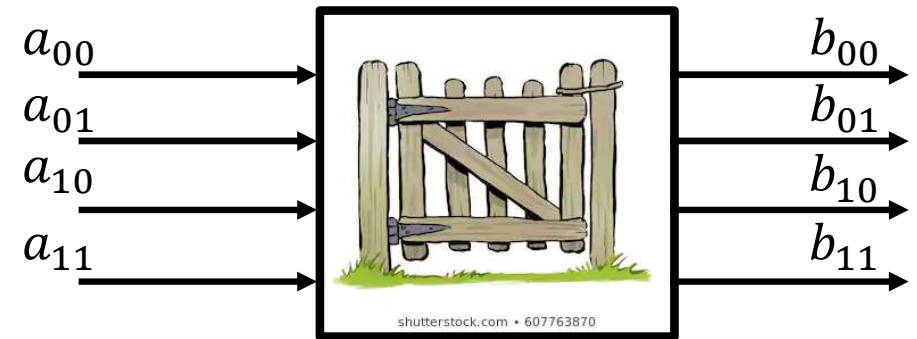
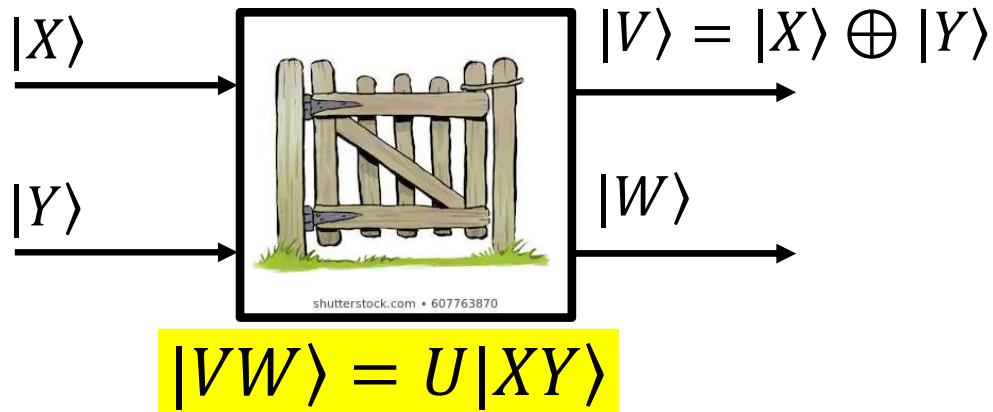
- Are o_1 and o_2 entangled?

CNOT IS AN ENTANGLING GATE!!

Lets try some simple gates

- So we know how to construct simple 2-qubit quantum gates
- So now, lets try to build them for the following 2-input Boolean operations
 - $X \oplus Y$
 - $X (AND) Y$

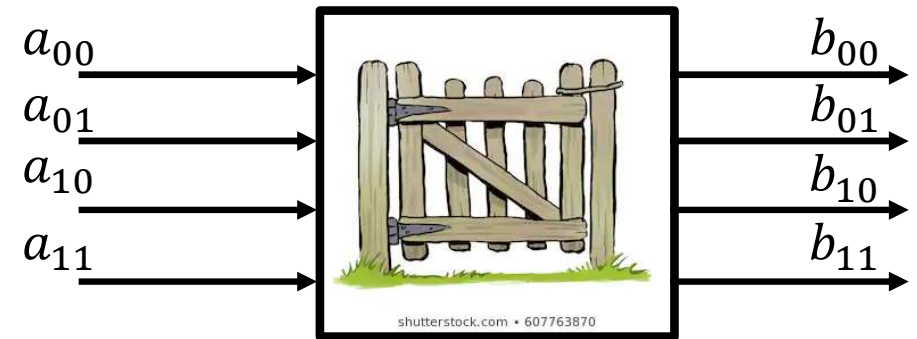
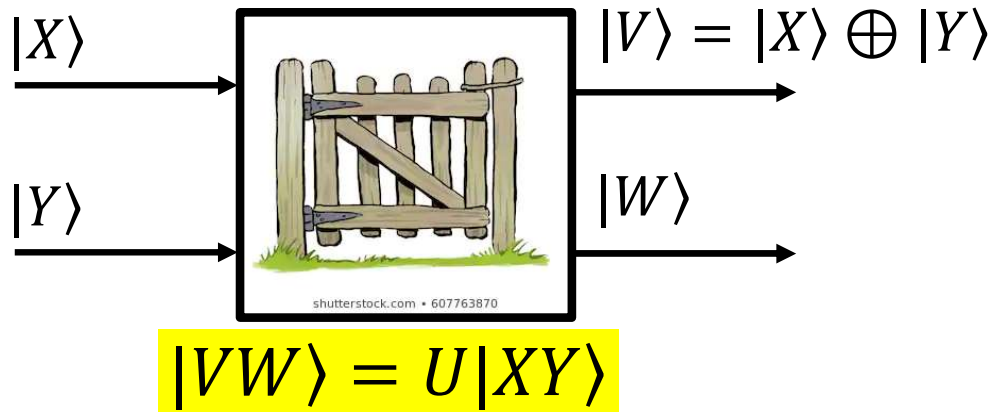
XOR



- We don't care what W is
- Create U such that the truth table to the right is produced

X	Y	V	W
0	0	0	
0	1	1	
1	0	1	
1	1	0	

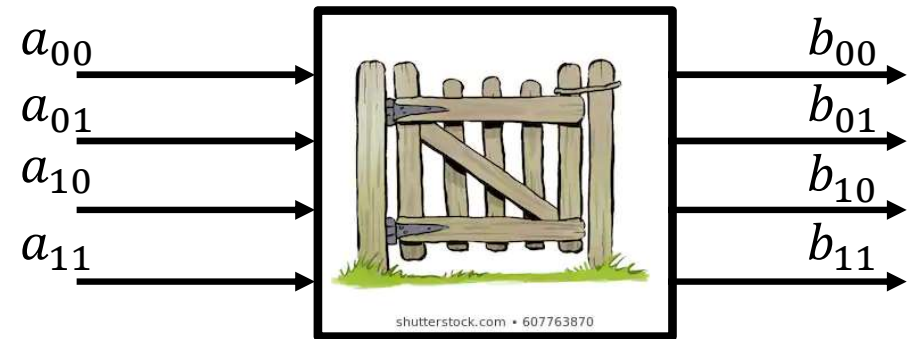
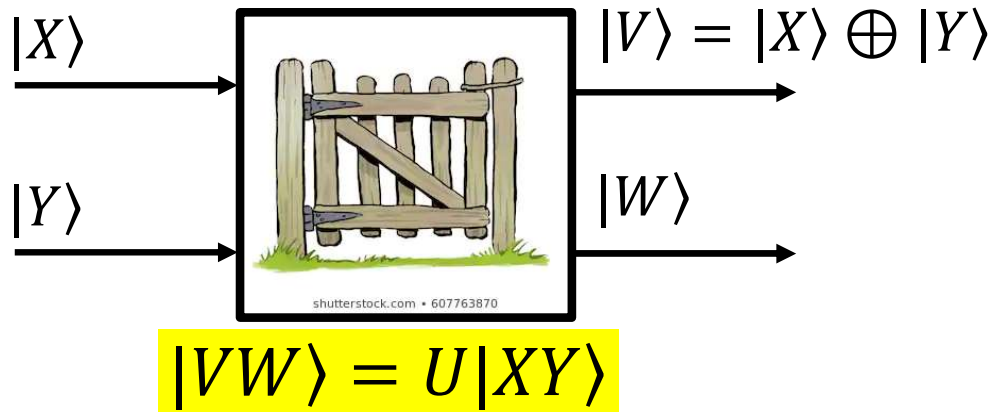
XOR



- $U|00\rangle = |0*\rangle$
- How do we select W to construct a U for this?

X	Y	V	W
0	0	0	
0	1	1	
1	0	1	
1	1	0	

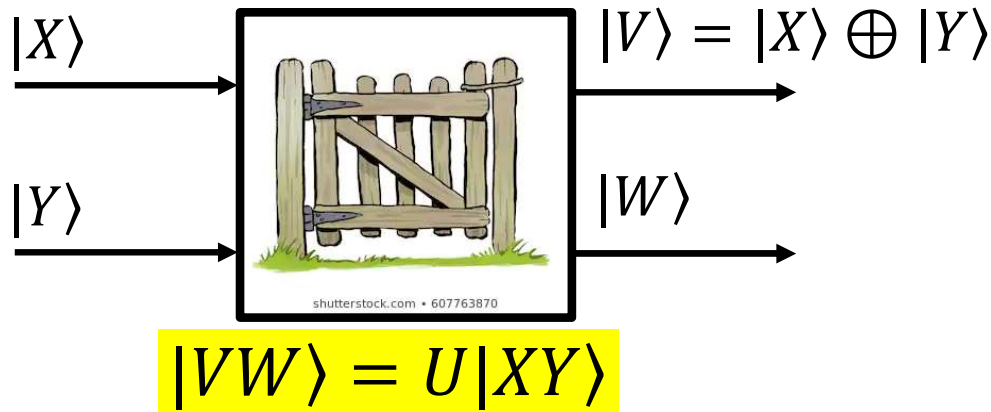
XOR



- $U|00\rangle = |0*\rangle$
- How do we select W to construct a U for this?
- Invertibility: Every orange row must be unique (or the function is not invertible)
 - Create W for this
- Design a U for the chosen W

X	Y	V	W
0	0	0	0
0	1	1	0
1	0	1	1
1	1	0	1

XOR



X	Y	V	W
0	0	0	0
0	1	1	0
1	0	1	1
1	1	0	1

- $U|00\rangle = |00\rangle, U|01\rangle = |10\rangle, U|10\rangle = |11\rangle, U|11\rangle = |01\rangle$

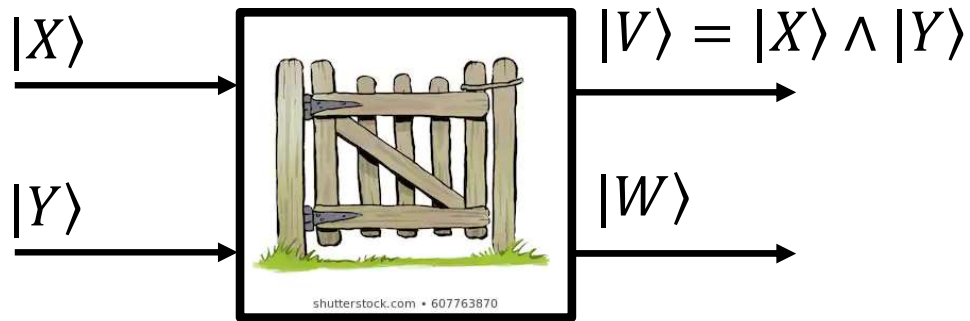
$$\begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ u_{10} & u_{11} & u_{12} & u_{13} \\ u_{20} & u_{21} & u_{22} & u_{23} \\ u_{30} & u_{31} & u_{23} & u_{33} \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ u_{10} & u_{11} & u_{12} & u_{13} \\ u_{20} & u_{21} & u_{22} & u_{23} \\ u_{30} & u_{31} & u_{23} & u_{33} \end{bmatrix} \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$$

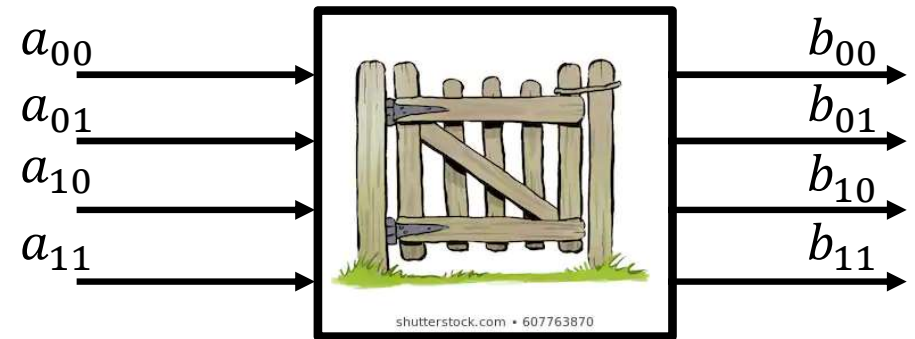
$$\begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix} = \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ u_{10} & u_{11} & u_{12} & u_{13} \\ u_{20} & u_{21} & u_{22} & u_{23} \\ u_{30} & u_{31} & u_{23} & u_{33} \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ u_{10} & u_{11} & u_{12} & u_{13} \\ u_{20} & u_{21} & u_{22} & u_{23} \\ u_{30} & u_{31} & u_{23} & u_{33} \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

AND



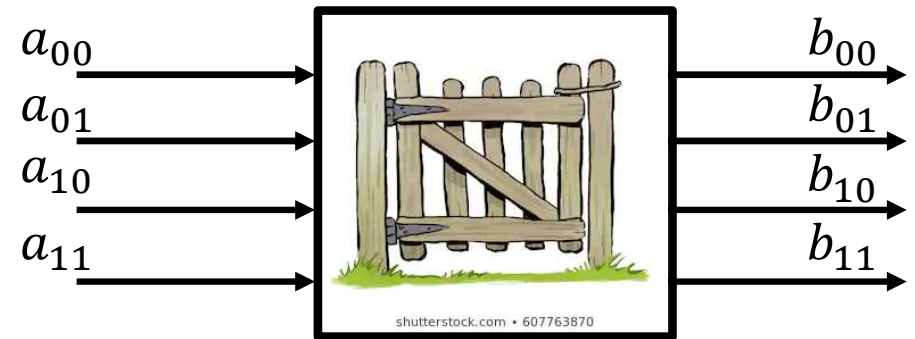
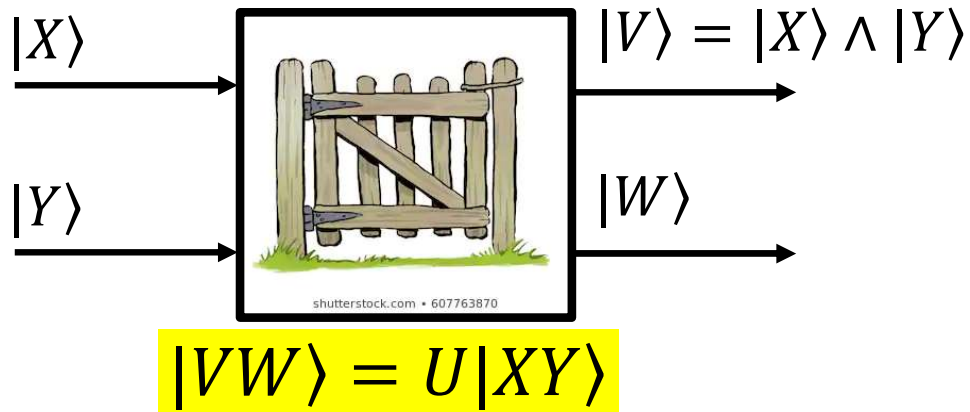
$$|VW\rangle = U|XY\rangle$$



- How do we select W to be able to construct a U for this table?

X	Y	V	W
0	0	0	*
0	1	0	*
1	0	0	*
1	1	1	*

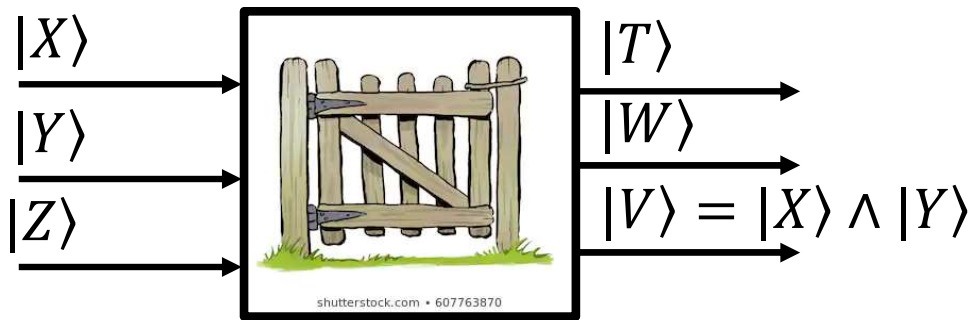
AND



- How do we select W to be able to construct a U for this table?
- You cannot!!!
 - Which other gates are similarly impossible to model?

X	Y	V	W
0	0	0	*
0	1	0	*
1	0	0	*
1	1	1	*

So how do you do an AND



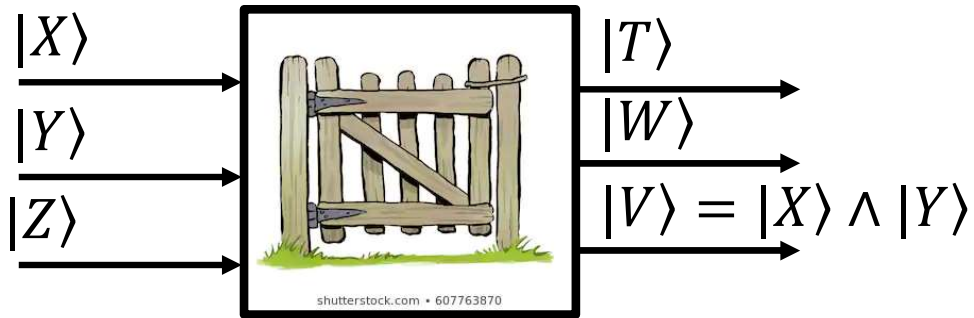
$$|VW\rangle = U|XY\rangle$$



- You add an *output qubit*
 - The fact is, the output is not manufactured from thin air
 - Its actually an entire qubit
- Does this help?

X	Y	Z	T	W	V
0	0	0	*	*	*
0	0	1	*	*	*
0	1	0	*	*	*
0	1	1	*	*	*
1	0	0	*	*	*
1	0	1	*	*	*
1	1	0	*	*	*
1	1	1	*	*	*

So how do you do an AND



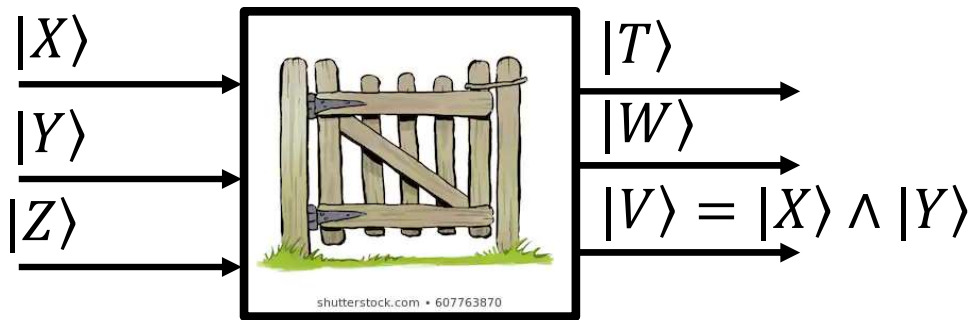
$$|VW\rangle = U|XY\rangle$$



- You can't in general
 - If you want V to give you the right input regardless of Z

X	Y	Z	T	W	V
0	0	0	*	*	0
0	0	1	*	*	0
0	1	0	*	*	0
0	1	1	*	*	0
1	0	0	*	*	0
1	0	1	*	*	0
1	1	0	*	*	1
1	1	1	*	*	1

So how do you do an AND



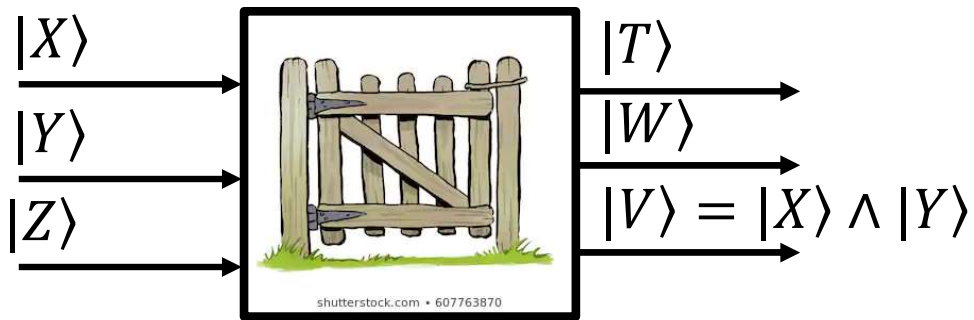
$$|VW\rangle = U|XY\rangle$$



- You can't in general
 - If you want V to give you the right input regardless of Z
- But it only has to work for *one value of Z*
 - This gives you a lot more freedom to “design” your transform

X	Y	Z	T	W	V
0	0	0	*	*	0
0	0	1	*	*	*
0	1	0	*	*	0
0	1	1	*	*	*
1	0	0	*	*	0
1	0	1	*	*	*
1	1	0	*	*	1
1	1	1	*	*	*

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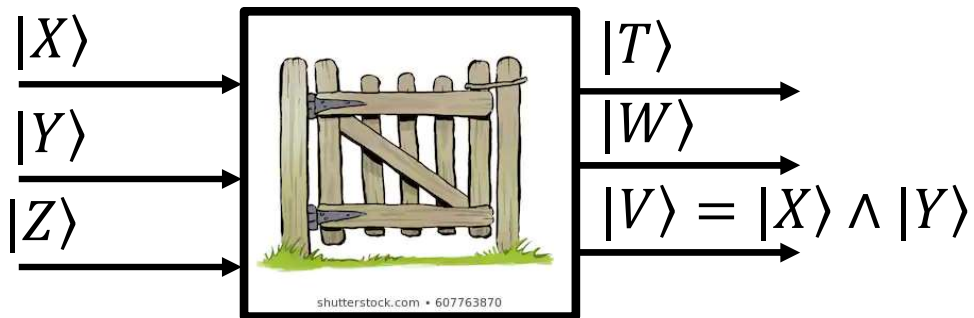
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X	Y	Z	T	W	V
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0	0	1	0	0	1
0	1	0	0	1	0
0	1	1	0	1	1
1	0	0	1	0	0
1	0	1	1	1	0
1	1	0	1	0	1
1	1	1	1	1	1

The CSWAP (or FREDKIN) Gate



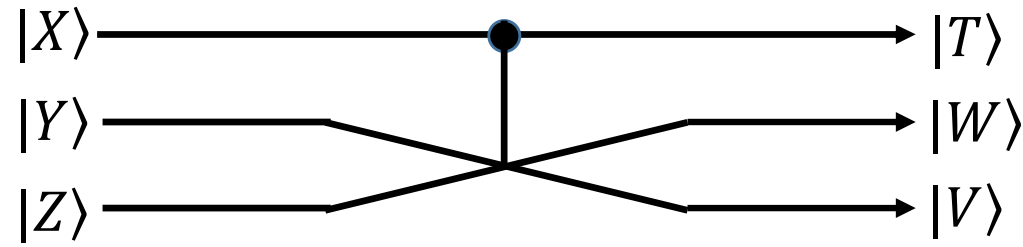
$$|VW\rangle = U|XY\rangle$$

$$U = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$



X	Y	Z	T	W	V
0	0	0	0	0	0
0	0	1	0	0	1
0	1	0	0	1	0
0	1	1	0	1	1
1	0	0	1	0	0
1	0	1	1	1	0
1	1	0	1	0	1
1	1	1	1	1	1

The CSWAP gate

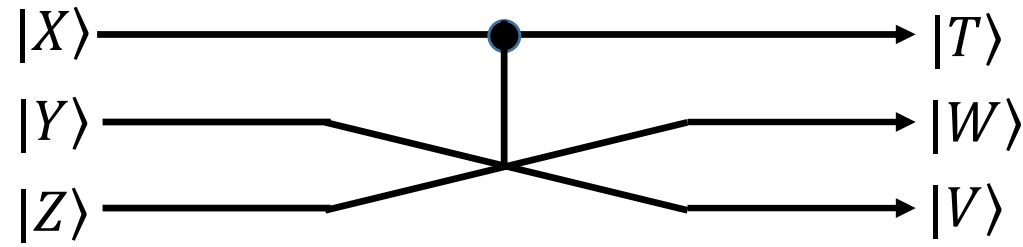


- The controlled swap gate
 - If the “control” X is 0, Y and Z go through in the same order
 - If X is 1, Y and Z swap
- Verify that when $Z = 0$, $V = X \text{ AND } Y$
- What other Boolean functions can you get by varying Z ?
 - And at what output variable?
- What would the output be for *superposed* states?

The CSWAP gate

Note: The CSWAP, combined with a NOT is a universal gate!!

Why??

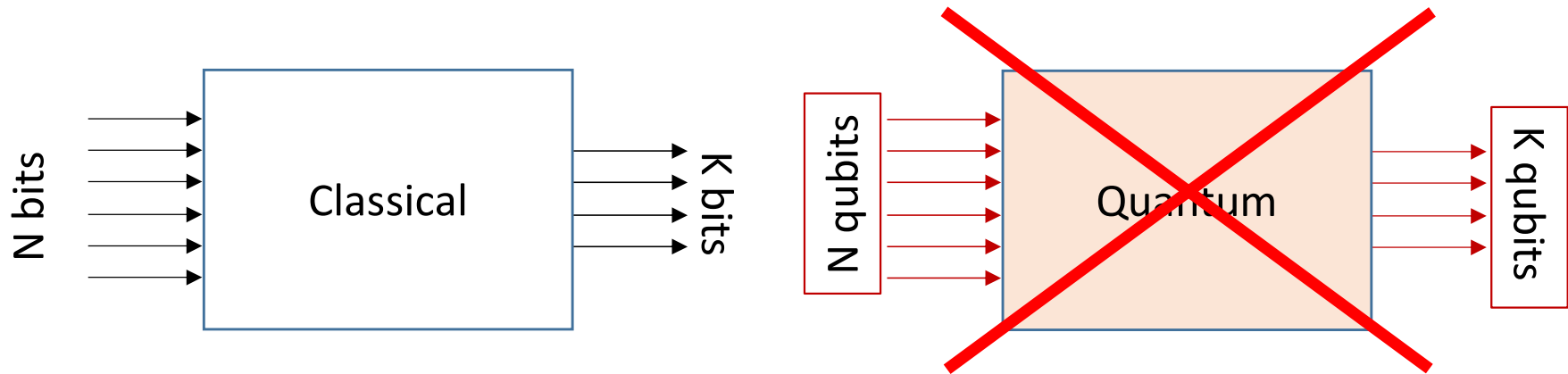


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Lesson for the day

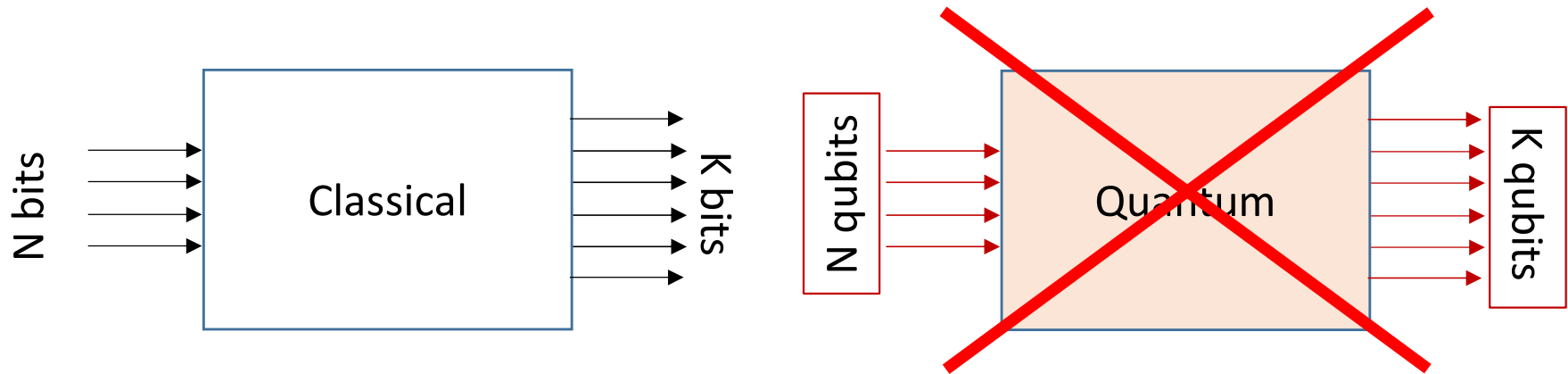
- You cannot simply emulate N-bit classical gates with an N-bit quantum circuit
- You will have to add extra qubits to hold the output
- And still more qubits to hold other necessary variables
 - AKA “junk”

Classical vs Quantum circuits



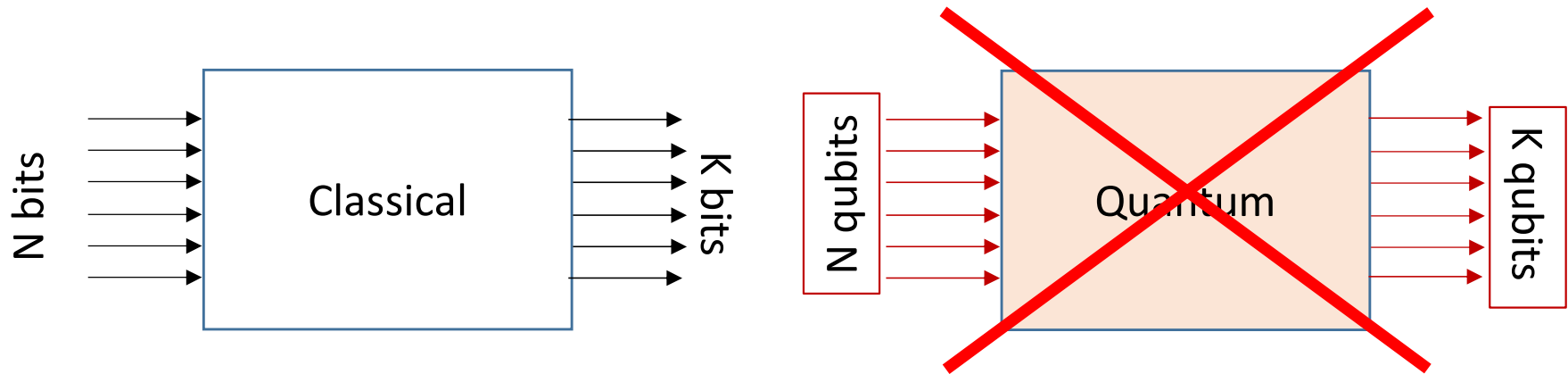
- Classical circuits:
 - N bits go in
 - K bits come out
- Quantum circuits:
 - Can often not directly emulate the classical circuit (with N input qubits and K output qubits)
 - For $K < N$, can *definitely* not emulate the classical circuit directly

Classical vs Quantum circuits



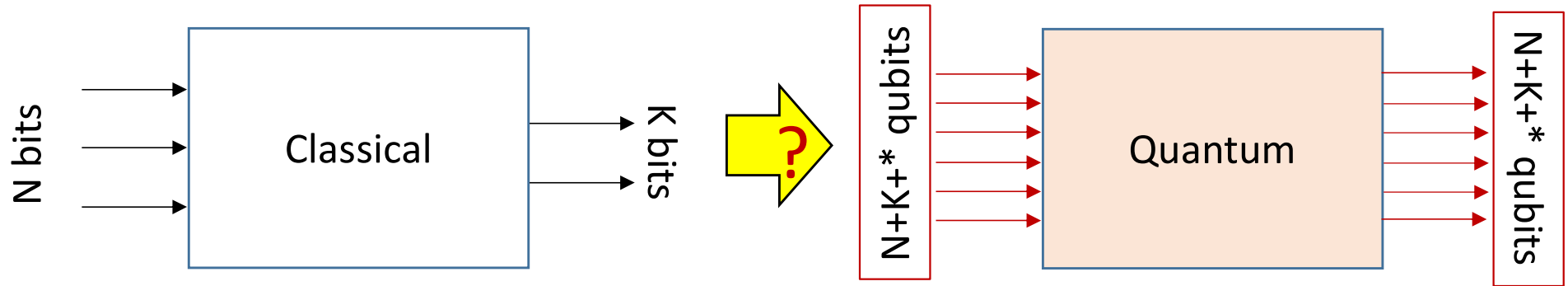
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Classical vs Quantum circuits



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 - Even for $K = N$, often need additional inputs and outputs

Classical vs Quantum circuits



- Classical circuits:

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- K bits come out

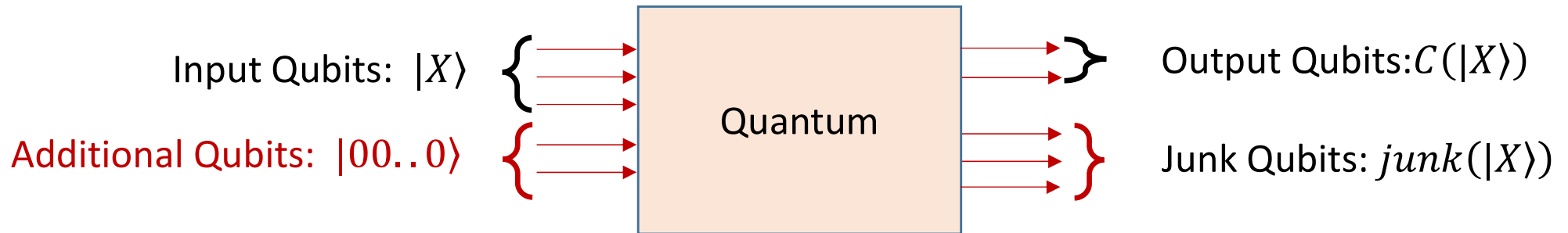
- What principles do we use to design them?

- Is there a *generic* method?
- And what must we watch out for?

- Quantum circuits:

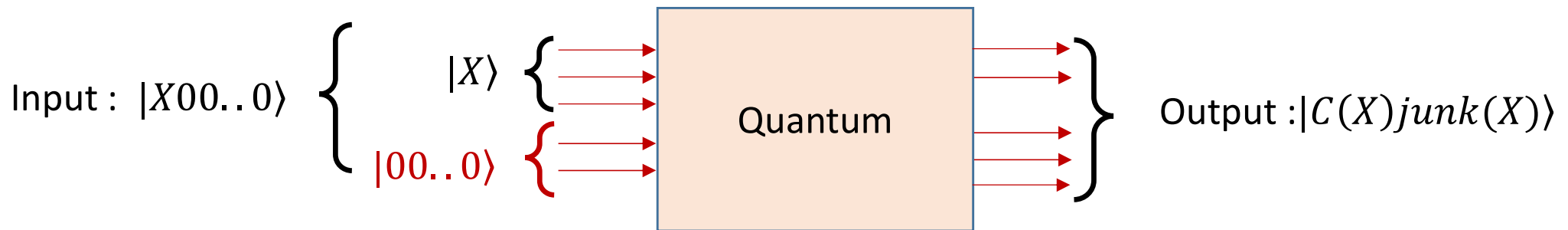
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Making a classical circuit quantum



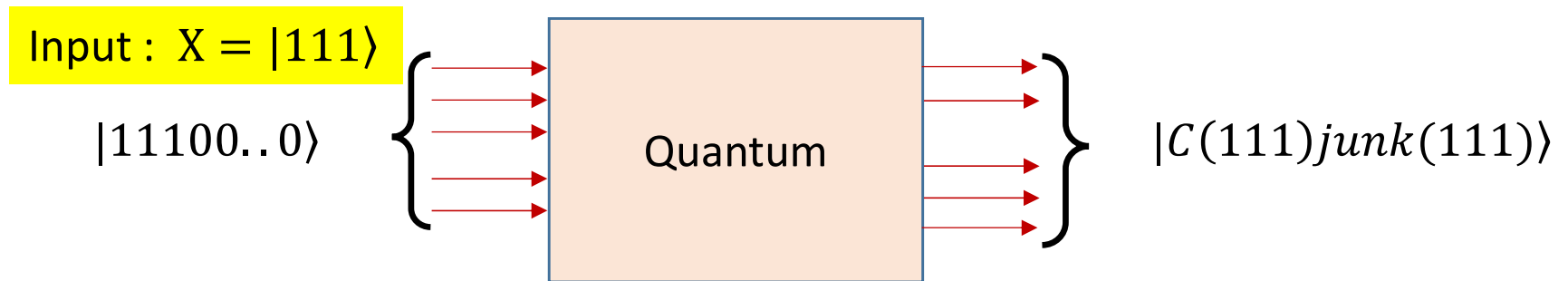
- First, make it reversible
 - Total number of input qubits = Total number of output qubits
- Output bits don't just emerge from the ether, they were always there
 - So, actual circuit has as input " X bits | Y bits"
 - The output is actually " $C(X)$ | $junk(X)$ "
 - The input bits too may get modified
 - No. of qubits in $junk(X)$ = no. of qubits in X
- But... Not so simple...

Making a classical circuit quantum



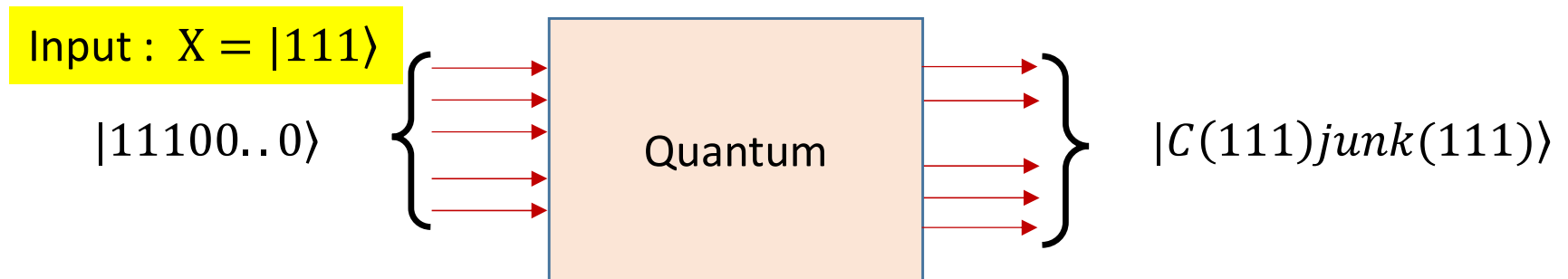
- The output is not $C(X) \otimes junk(X)$
 - It is actually $|C(X)junk(X)\rangle$
 - I.e. for any given input X you will not be able to obtain all possible combinations of $C(X)$ and all 2^N possible values of $junk(X)$ simply by manipulating the *input* values of the output bits Y
 - Why is this the case?

Making a classical circuit quantum



- The output is not $C(X) \otimes junk(X)$
 - It is actually $|C(X)junk(X)\rangle$
- When you fix X , the values of $junk(X)$ are restricted
 - Only K bits available to manipulate N bits
- In other words, when you fix $C(X)$, the values of $junk(X)$ are restricted
 - The target output bits and the junk output bits are *entangled!!!*

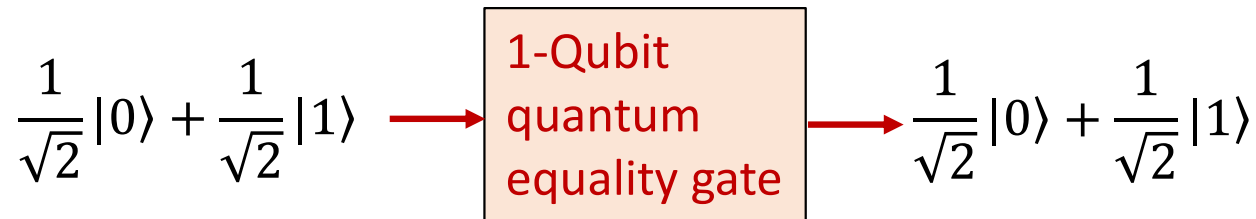
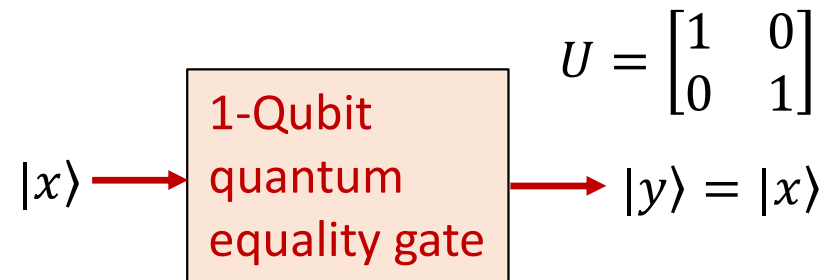
Making a classical circuit quantum



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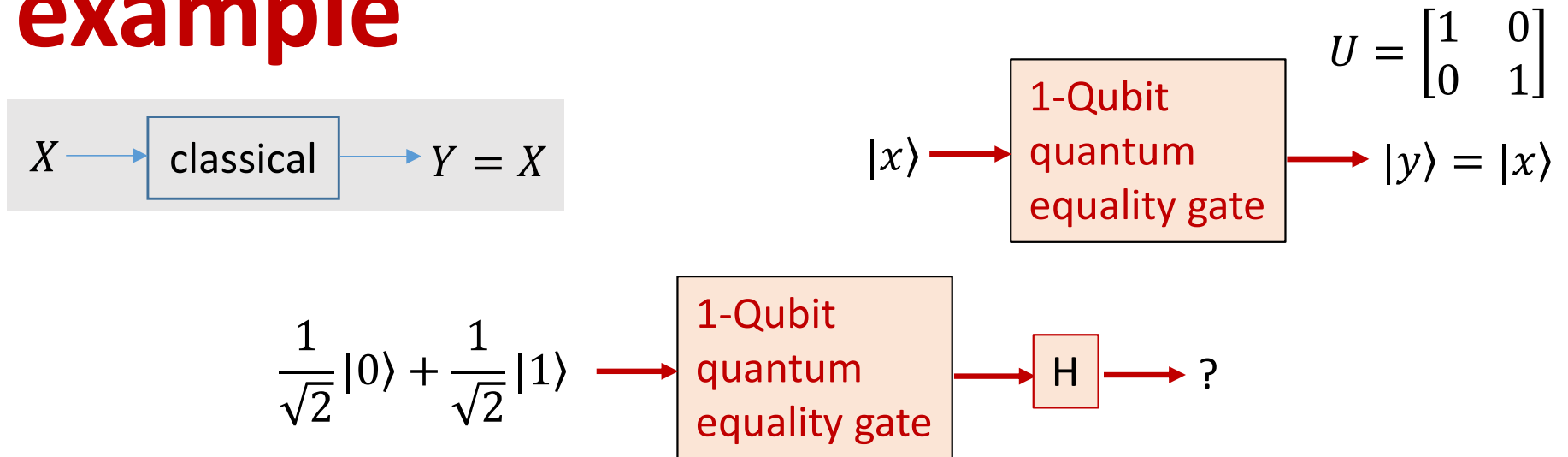
Why is this a problem?

The trouble with junk – an example



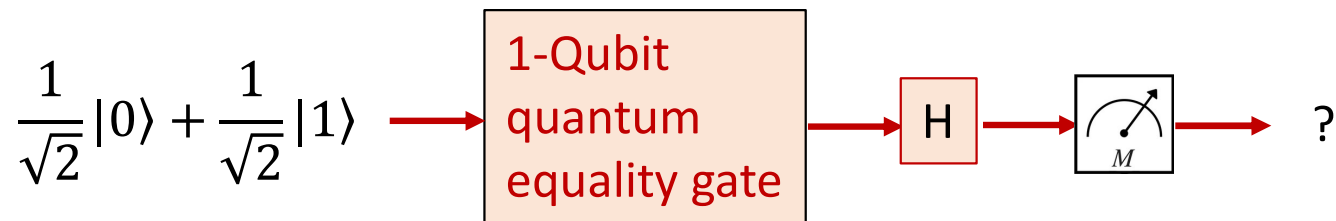
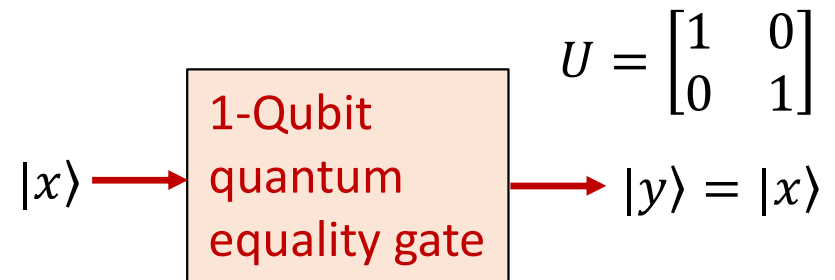
- Consider a 1-bit equality gate
 - Classically $y = x$
- This can in fact be implemented using a 1-qubit quantum gate with $U = I$
 - More generally, if $x = a_0|0\rangle + a_1|1\rangle$, then the output $y = a_0|0\rangle + a_1|1\rangle$
- We input the sign basis $|x\rangle = |+\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$
 - Output is $|y\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$

The trouble with junk – an example



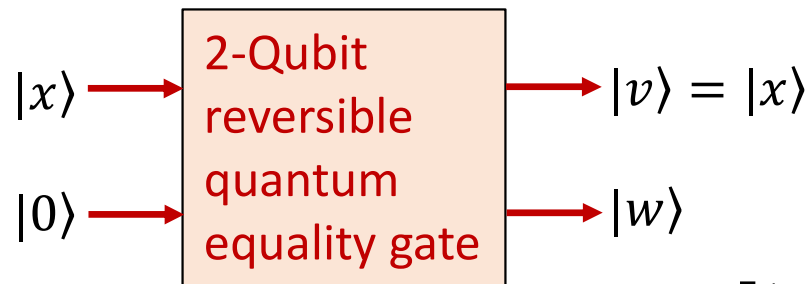
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 - Output is $|y\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$
- Now we put a Hadamard gate at the output
 - What is the output?
 - What will we get if we measure it?

The trouble with junk – an example



- Consider a 1-bit equality gate
 - Classically $y = x$
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The reversible 2-qubit equality gate

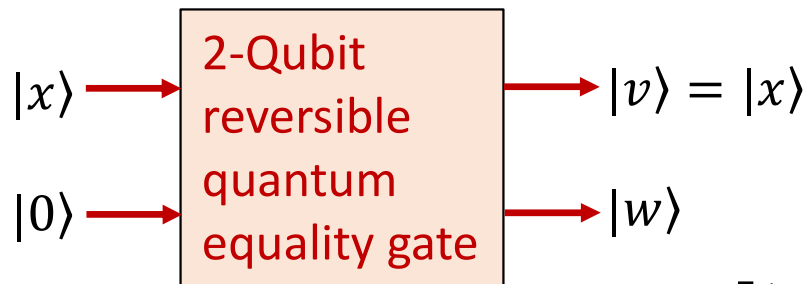


X	Y	V	W
0	0	0	0
0	1	0	1
1	0	1	1
1	1	1	0

$$U = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

- Now we convert it to a reversible gate using the earlier formula
 - Add a second input and second output
 - We construct the truth table above:
- The equality gate is obtained by setting the second input to 0
 - For $|y\rangle = |0\rangle$, the output $|v\rangle$ is $|v\rangle = |x\rangle$

The reversible 2-qubit equality gate



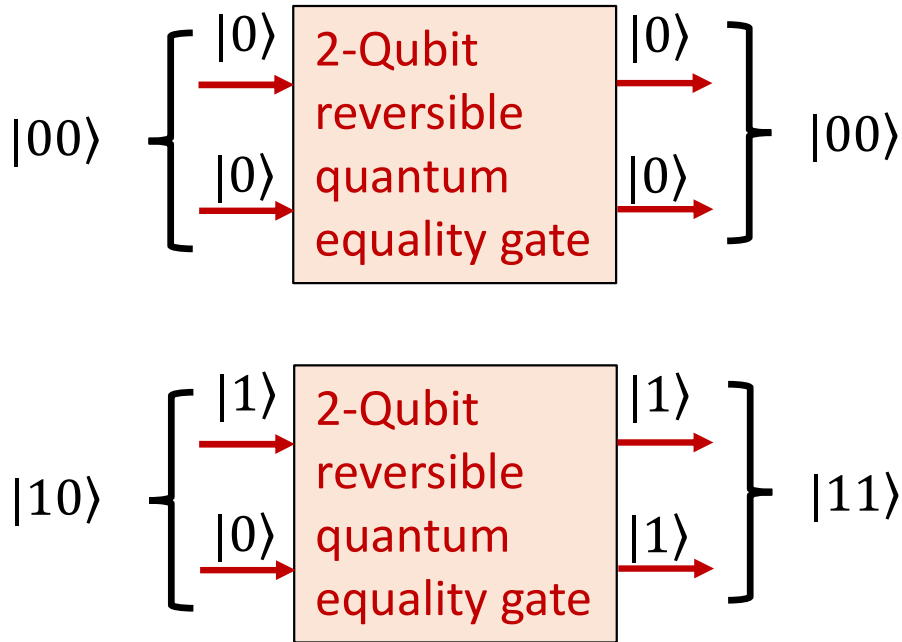
X	Y	V	W
0	0	0	0
0	1	0	1
1	0	1	1
1	1	1	0

$$U = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Note that this is not necessary for this gate; this exercise is meant for illustrating the problem with junk

- Now we convert it to a reversible gate using the earlier formula
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 - We construct the truth table above:
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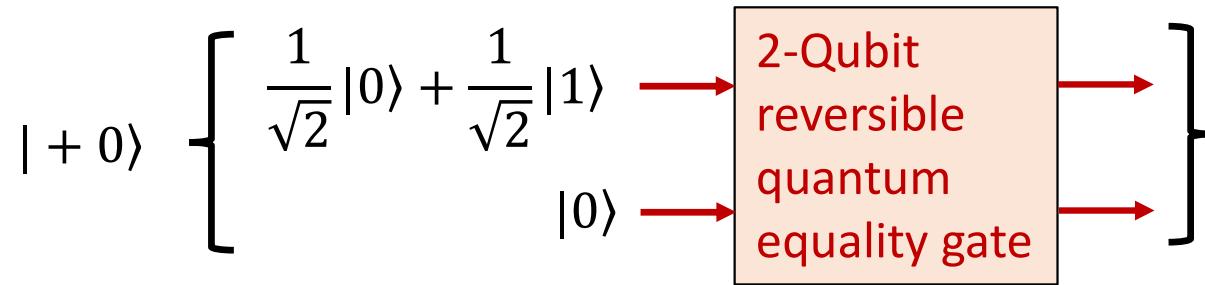
Looking closer



X	Y	V	W
0	0	0	0
0	1	0	1
1	0	1	1
1	1	1	0

- The equality gate is obtained by setting the second input to 0
- For $|y\rangle = |0\rangle$, the output $|v\rangle$ is *apparently* $|v\rangle = |x\rangle$
- The *actual* complete output is,
 - for $|xy\rangle = |00\rangle$, $\Rightarrow |vw\rangle = |00\rangle$
 - for $|xy\rangle = |10\rangle$, $\Rightarrow |vw\rangle = |11\rangle$

Lets input a + basis

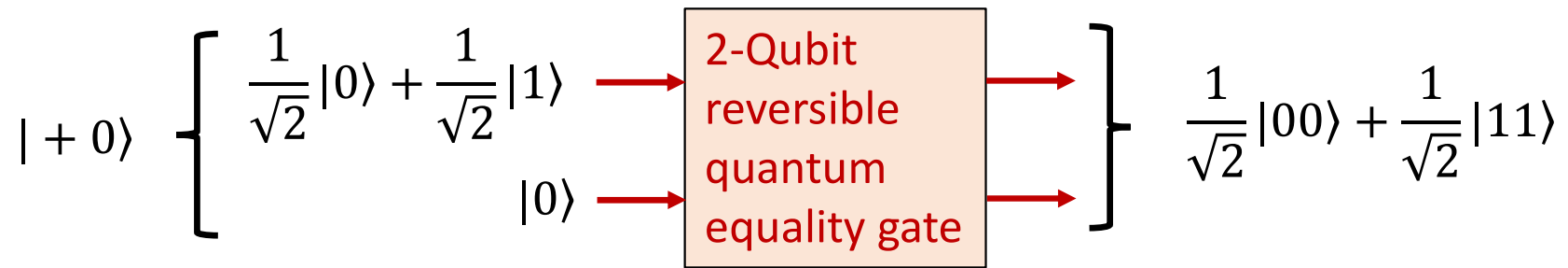


- The input is:

$$|x0\rangle = |+0\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|10\rangle$$

- The complete output is:

Lets input a + basis



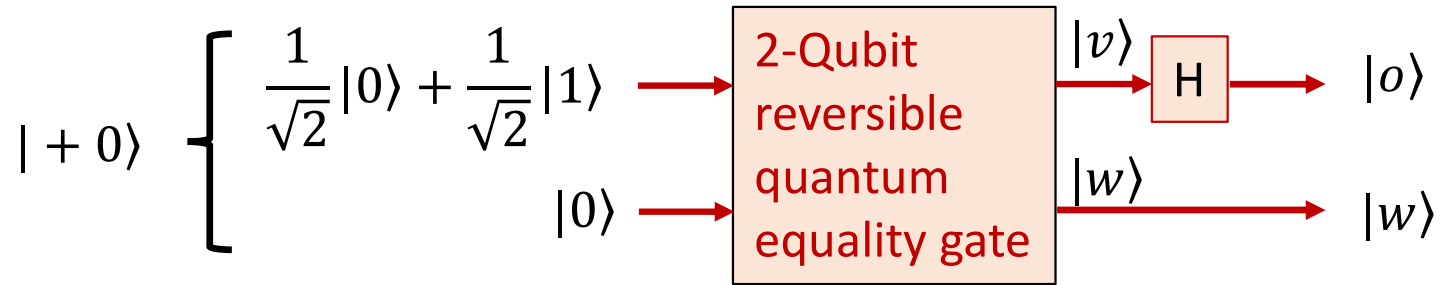
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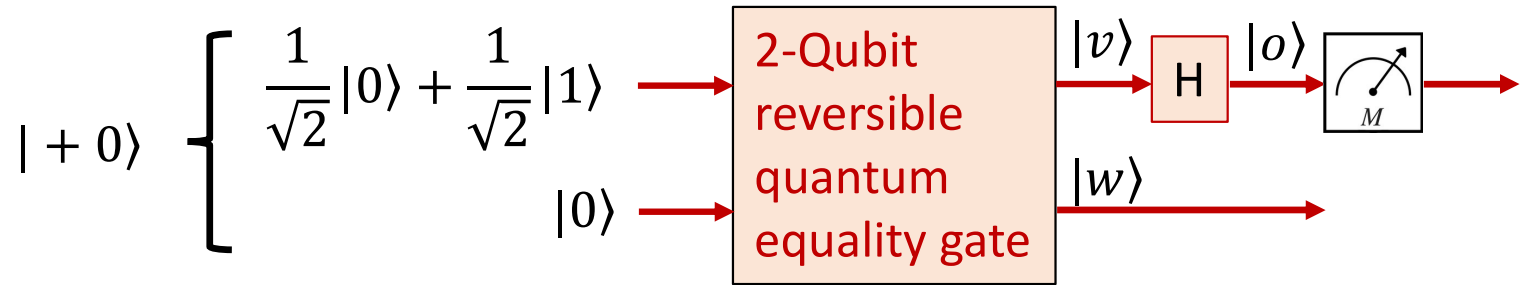
$$|vw\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$$

Collapsing the equality of the reversible 2-qubit equality gate



- Now add a Hadamard to the first qubit
- Note that this is effectively the same situation as we had with the one-qubit gate
- But is the output the same as for the 1-qubit gate?

The trouble with junk



- The actual output before the Hadamard is

$$|vw\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$$

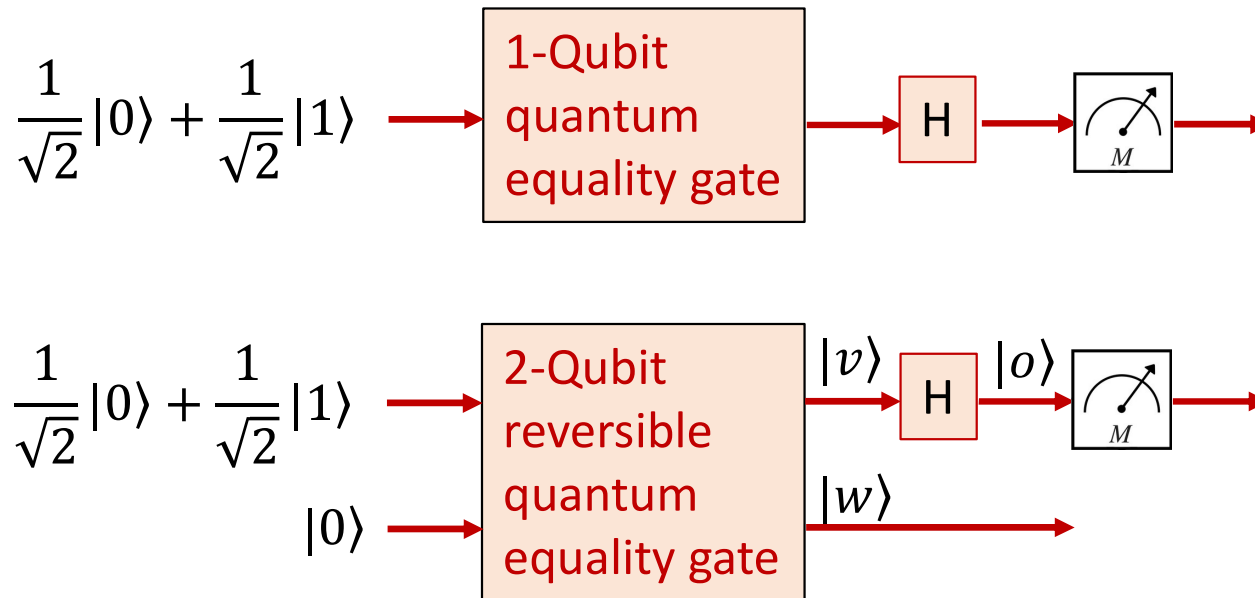
- The output after the Hadamard is:

$$|ov\rangle = \frac{1}{\sqrt{2}} \left| \left(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle \right) 0 \right\rangle + \frac{1}{\sqrt{2}} \left| \left(\frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle \right) 1 \right\rangle$$

$$|ov\rangle = \frac{1}{2}(|00\rangle + |10\rangle + |01\rangle - |11\rangle)$$

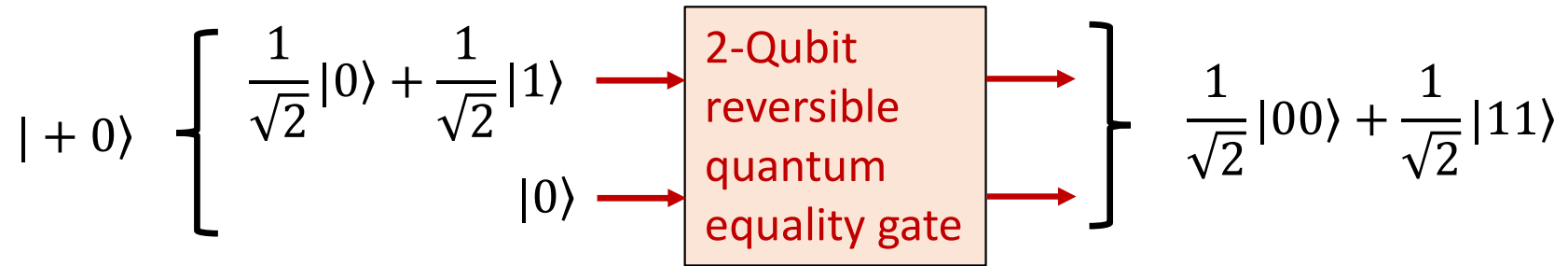
- The output is no longer deterministic. In fact the probability of measuring a 1 on the first bit is (what?)

The trouble with junk



- Comparison:
 - Using the 1-qubit gate the measured output is $|1\rangle$
 - Using the 2-qubit gate, the probability of measuring $|1\rangle$ is 0.5
- Simply having the junk bit destroyed our equality gate!!

Lets input a + basis

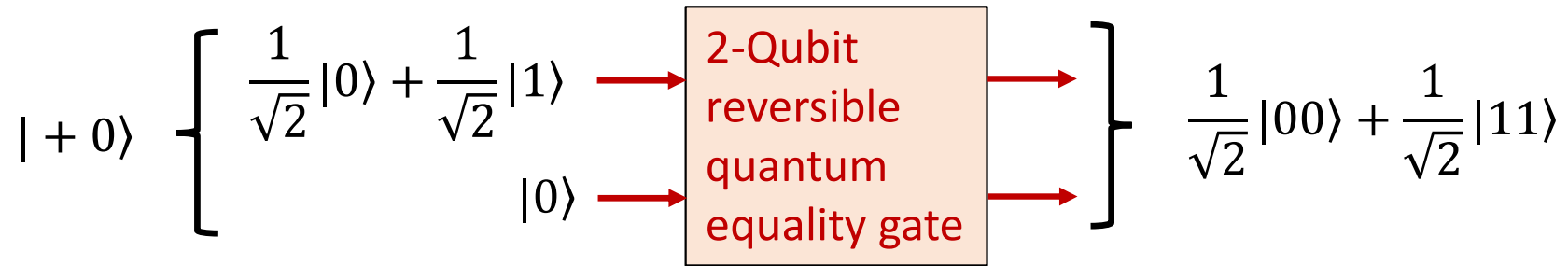


- The output is:

$$|vw\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$$

- If you measure w and get a 0, what is v ?
- If you measure w and get a 1, what is v ?

Lets input a + basis



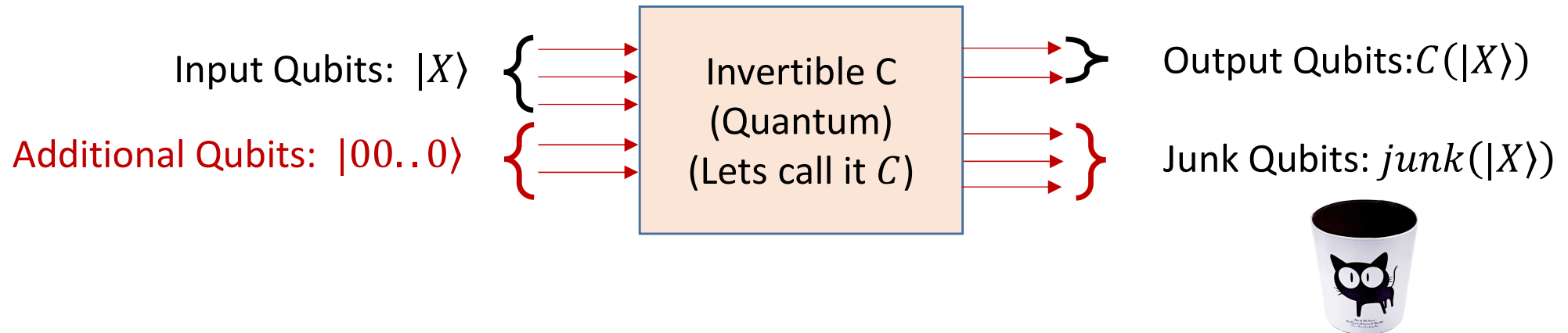
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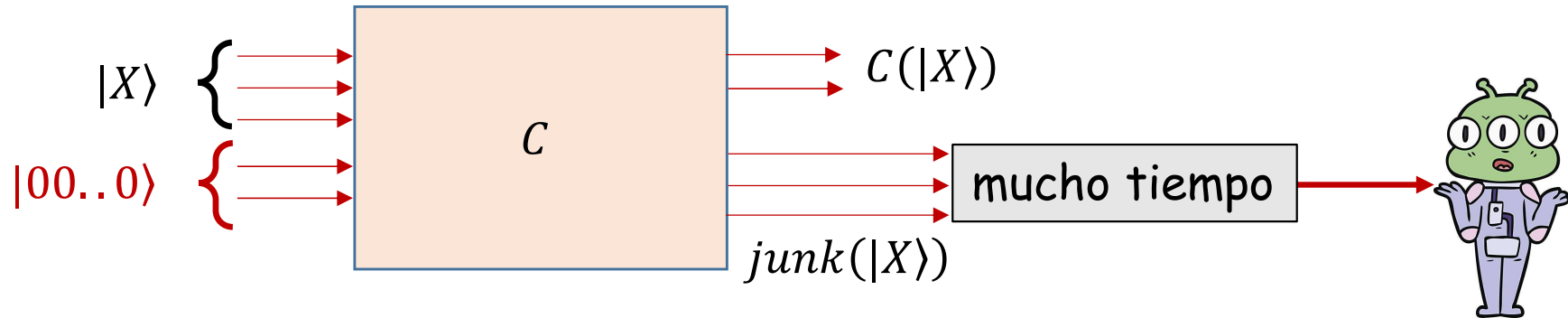
What we measure
on the junk bit
affects what you
measure on output

So why don't we just “junk” our junk bits?



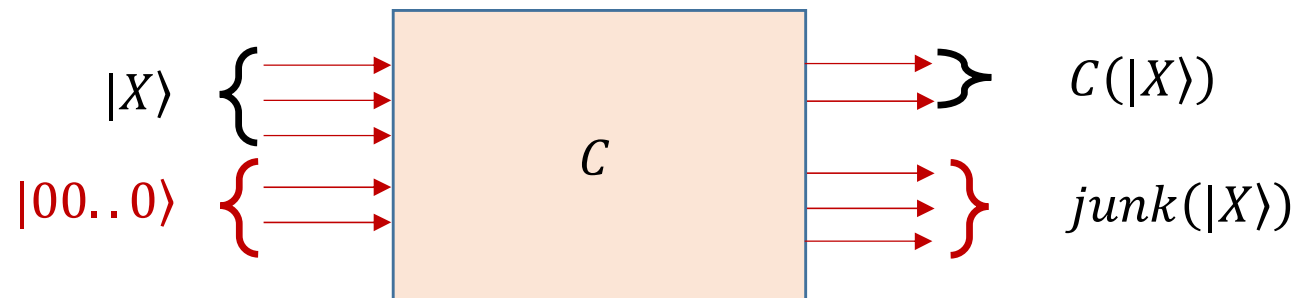
- Account for them in our arithmetic and ignore them...
- You can't

Remember the time machine?



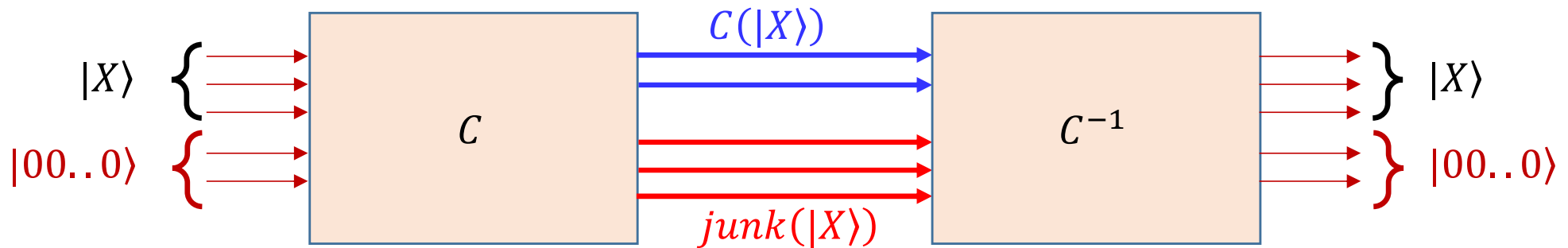
- Someone, somewhere, somewhen may measure the junk bits
 - Today or in the year 20,000,000AD
 - This someone could simply be nature
- What they measure will influence your measurement today
 - You cannot trust the output of your computation
 - You cannot simply assume the junk will never be measured, and you cannot assume what it will be measured as

So how do we deal with the junk?



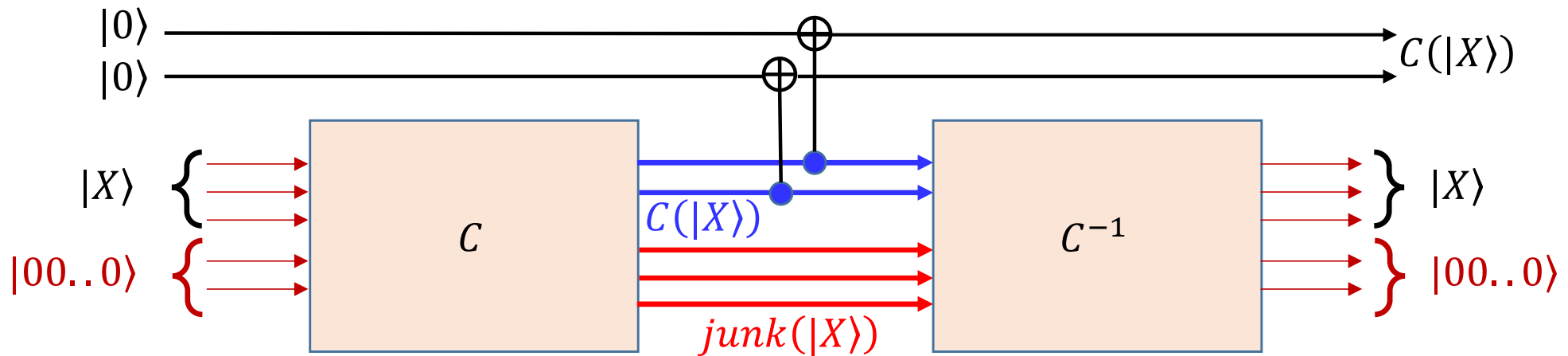
- The junk bits cannot just be discarded.
 - So how do we handle them?
- Desideratum : They must be “disentangled” from the actual output, somehow
- Hint: The circuit is invertible...

Eliminating the junk



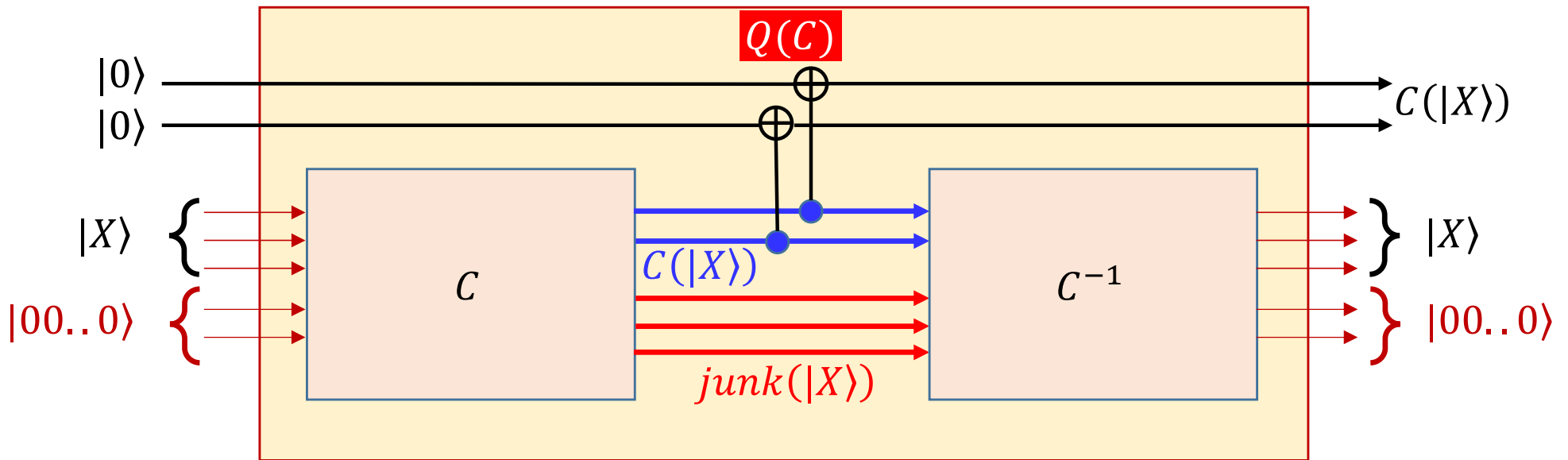
- If you connect the inverse of the circuit to the circuit, you “disentangle” the junk bits
 - Which will return to the value $|00 \dots 0\rangle$
- But now, we’ve lost the target output $C(|X\rangle)$

Retaining the output



- Have a second set of output qubits which are CNOTted with $C(|X\rangle)$
 - The output qubits are initialized to $|0\rangle$
 - $|0\rangle \text{ CNOT } C(X) = C(X)$
- The output is captured
 - The input is retained
 - The Junk bits are disentangled

The full circuit



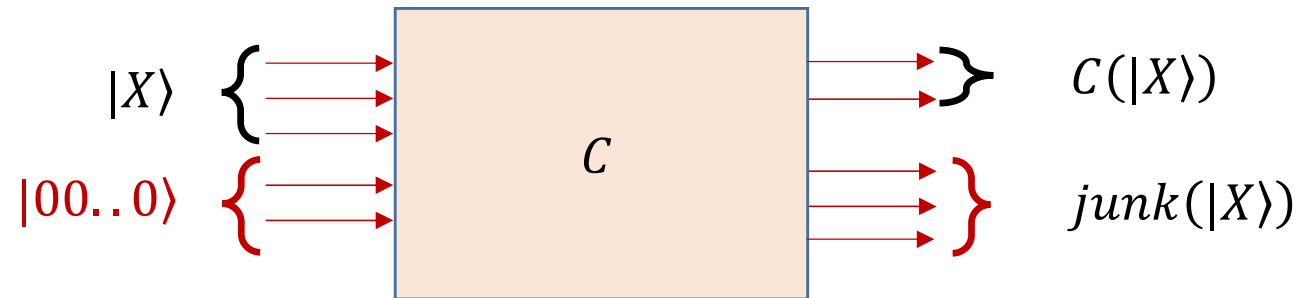
- Input comprises
 - $|X\rangle$
 - output bits initialized to $|0\rangle$
 - and a bunch of auxiliary $|0\rangle$ s needed for computation
- Output is
 - $C(|X\rangle)$,
 - $|X\rangle$,
 - and a bunch of $|0\rangle$ s,
- With junk disentangled

$|X\rangle$ remains entangled with $C(|X\rangle)$, as it must be

The full process: Step 1

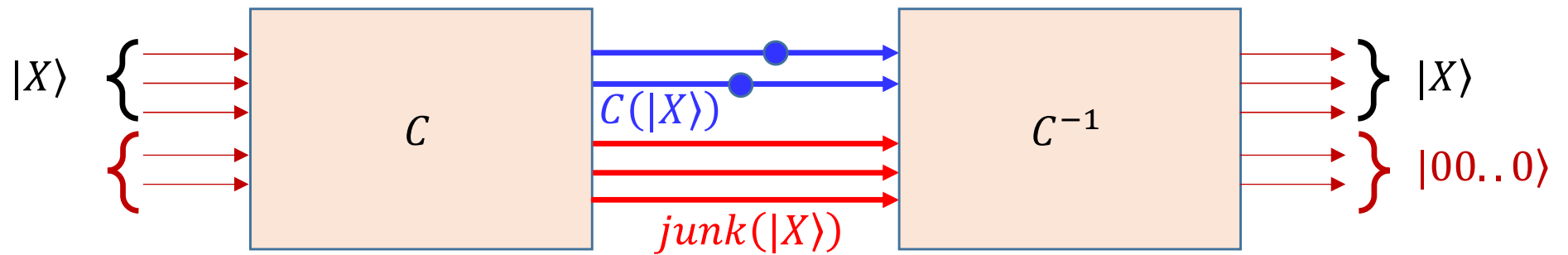
- Step 1a: Using truth tables
 - Compose a truth table for the function
 - Will require new output bits
 - Compose the transform for the table
 - As a (minimal) tensor product of universal set of quantum gates
 - Will generally require new junk bits
- Step 1alternative: Compose the circuit using quantum variants of Boolean gates
 - Construct quantum circuit using the gates
 - Will require output and junk bits

The output of step 1



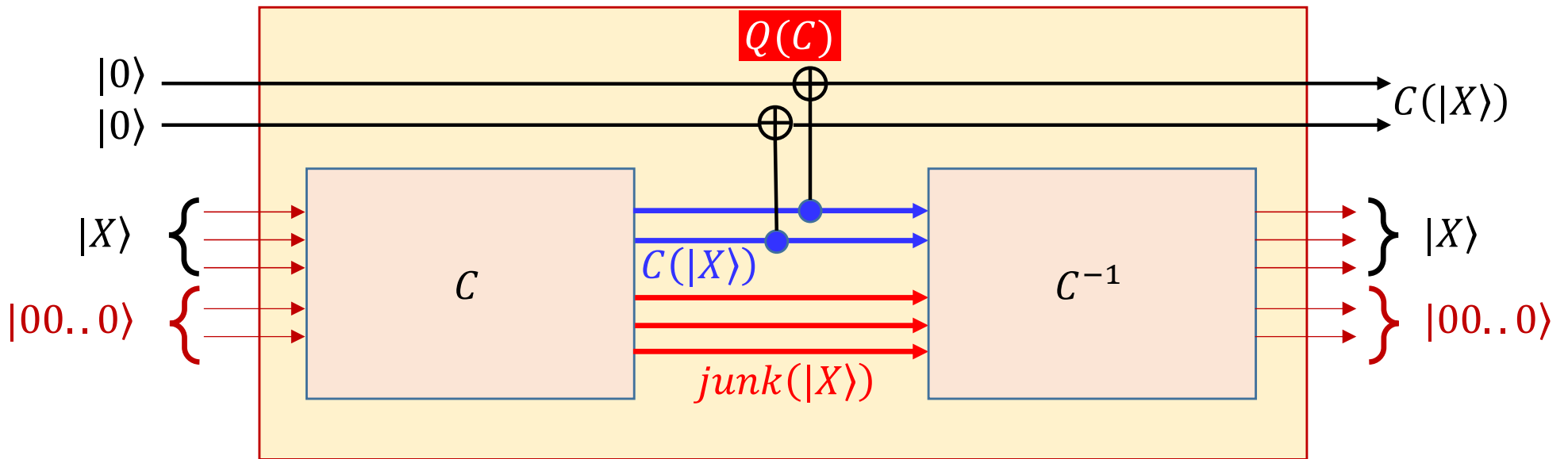
- A circuit that takes in input bits, output bits and junk bits
- And outputs the target output, plus a number of potentially entangled junk bits

Step 2:



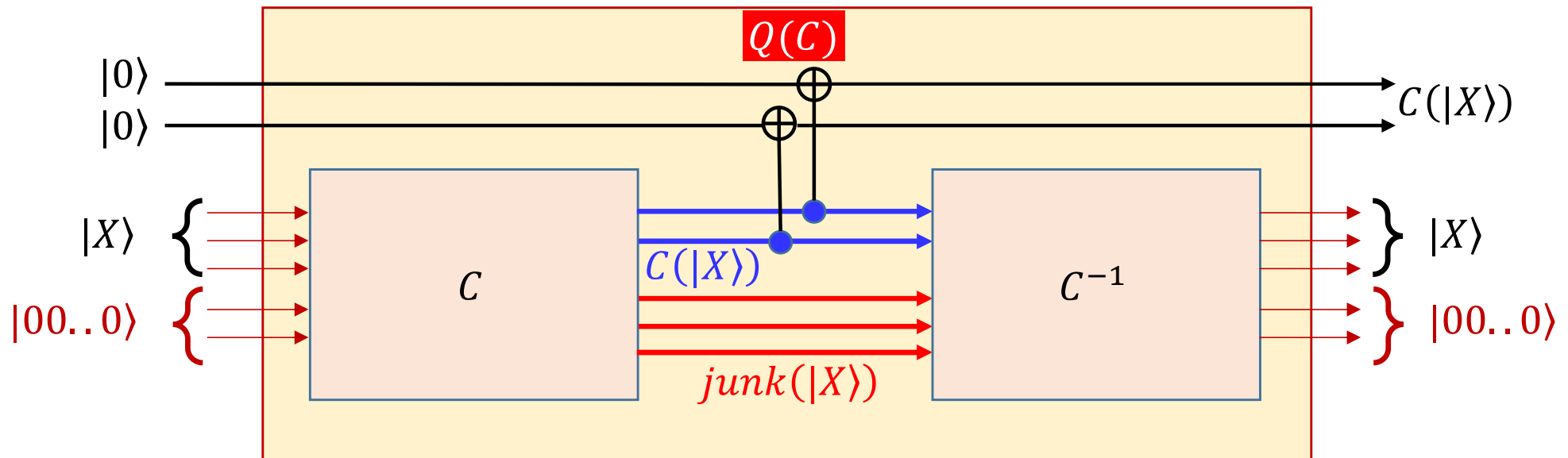
- Couple the output of Step 1 with its own inverse

Step 3:



- Add the actual output qubits, which are CNOTted with the $C(|X\rangle)$ computed by the inner circuit

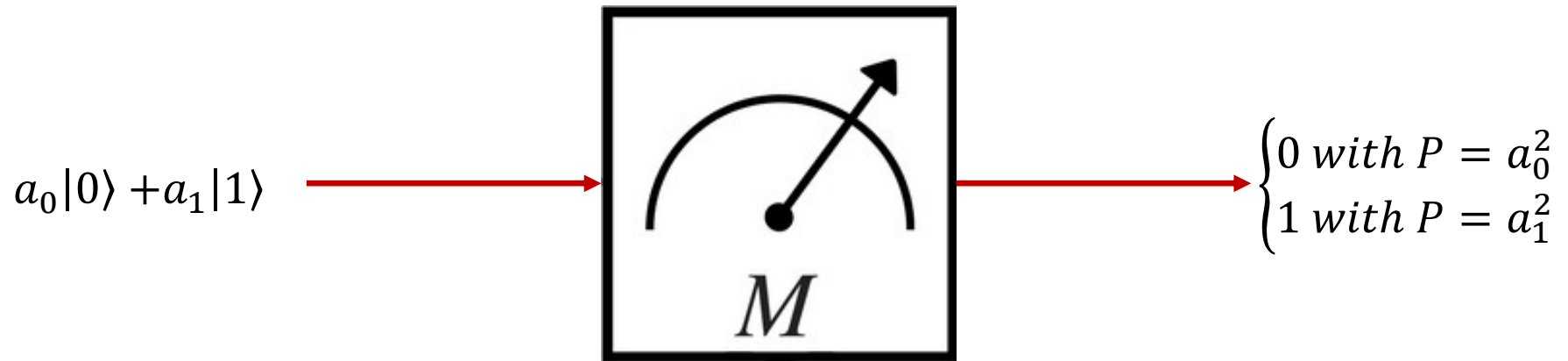
The full circuit



What are the universal quantum set of gates?

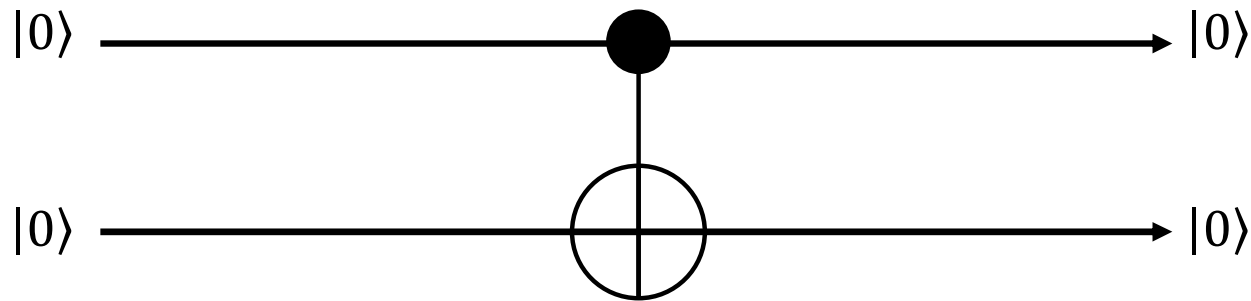
- Universal quantum set of gates:
 - CNOT
 - X
 - H
 - Z
 - $\frac{\pi}{8}$ rotation
- Any function can be computed using just these gates

A note on measurement...



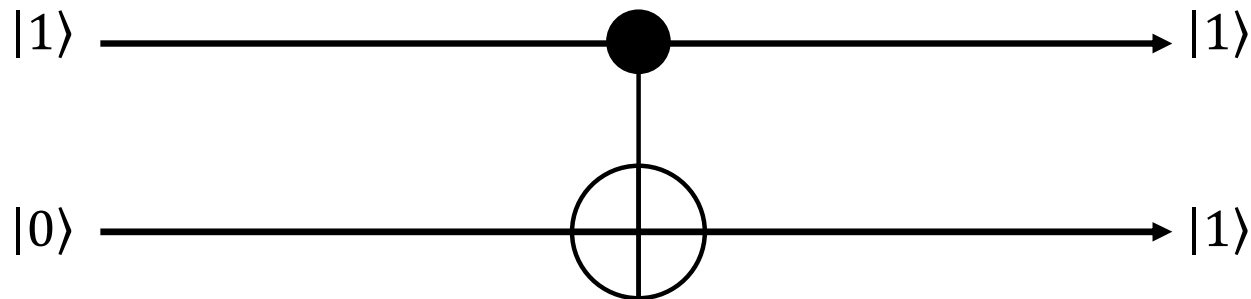
- What is this crazy thing called measurement?
- We have a qubit $|x\rangle = a_0|0\rangle + a_1|1\rangle$
- We want to run some physical operation on it such that the outcome is 0 with $P = a_0^2$ and 1 with $P = a_1^2$
- What might such a process look like?

The CNOT with $|0\rangle$ creates a Bell State



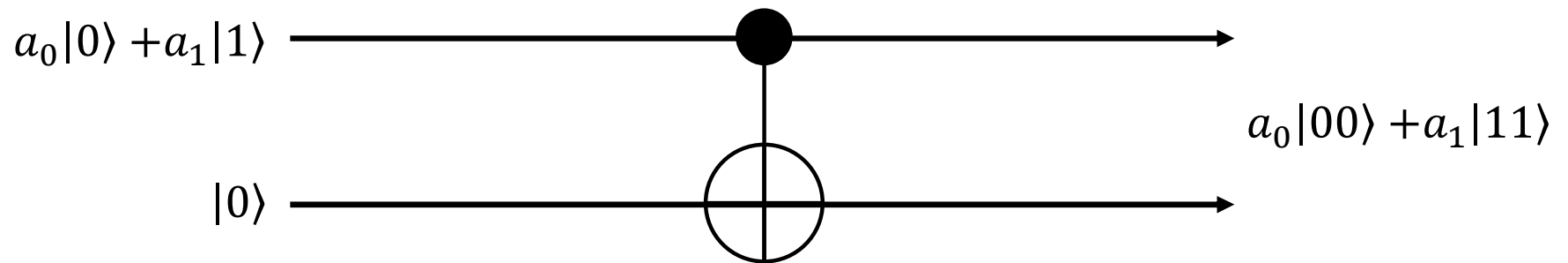
- The target input is $|x\rangle = a_0|0\rangle + a_1|1\rangle$
- The output is $|y\rangle = a_0|00\rangle + a_1|11\rangle$

The CNOT with $|0\rangle$ creates a Bell State



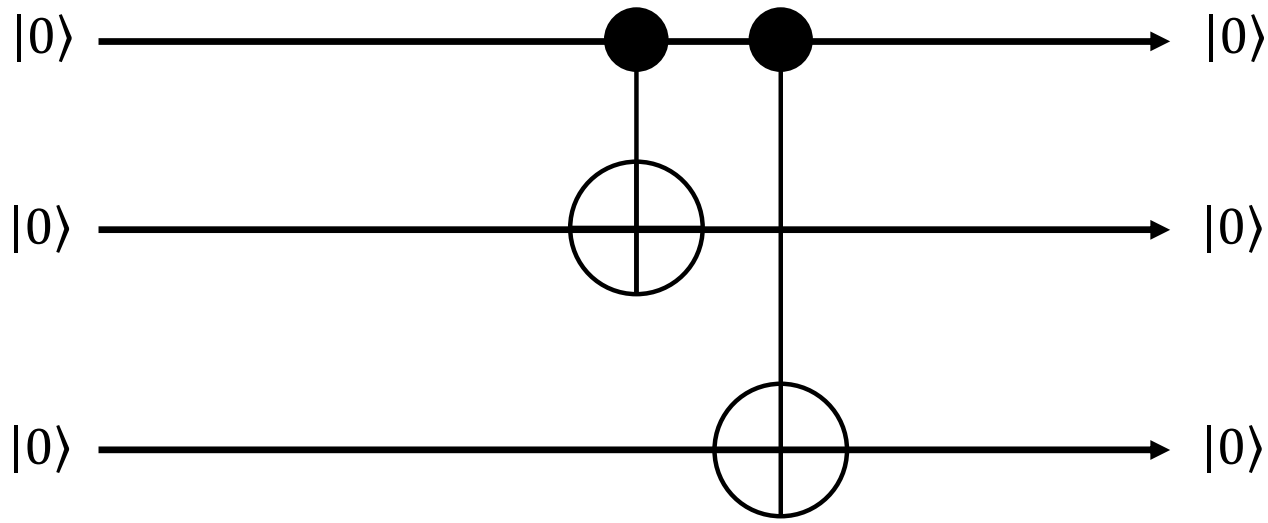
- The target input is $|x\rangle = a_0|0\rangle + a_1|1\rangle$
- The output is $|y\rangle = a_0|00\rangle + a_1|11\rangle$

The CNOT with $|0\rangle$ creates a Bell State



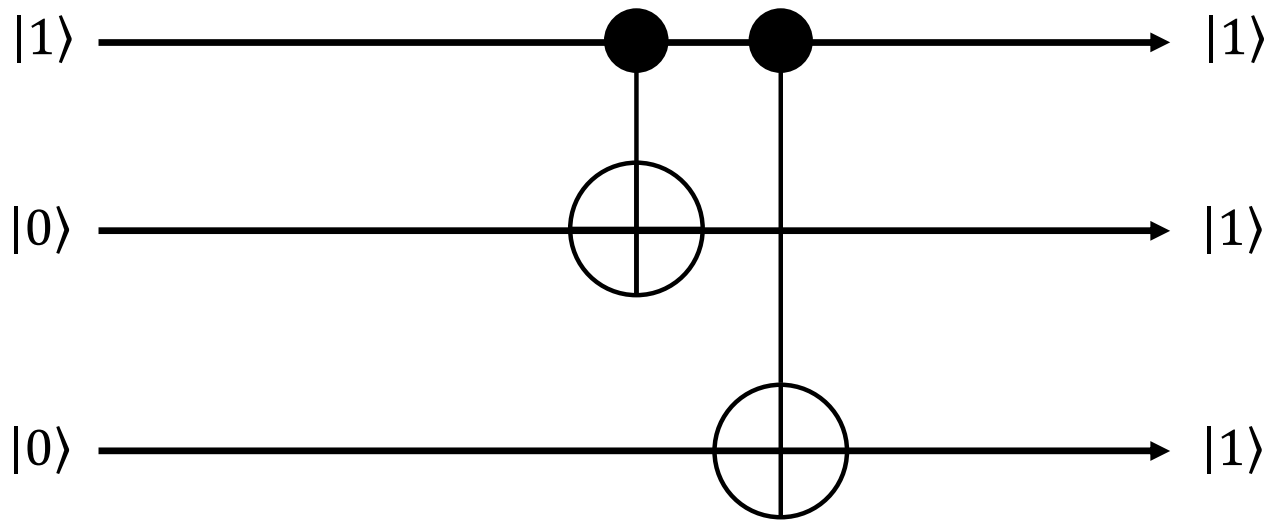
- The target input is $|x\rangle = a_0|0\rangle + a_1|1\rangle$
- The output is $|y\rangle = a_0|00\rangle + a_1|11\rangle$

Adding another CNOT



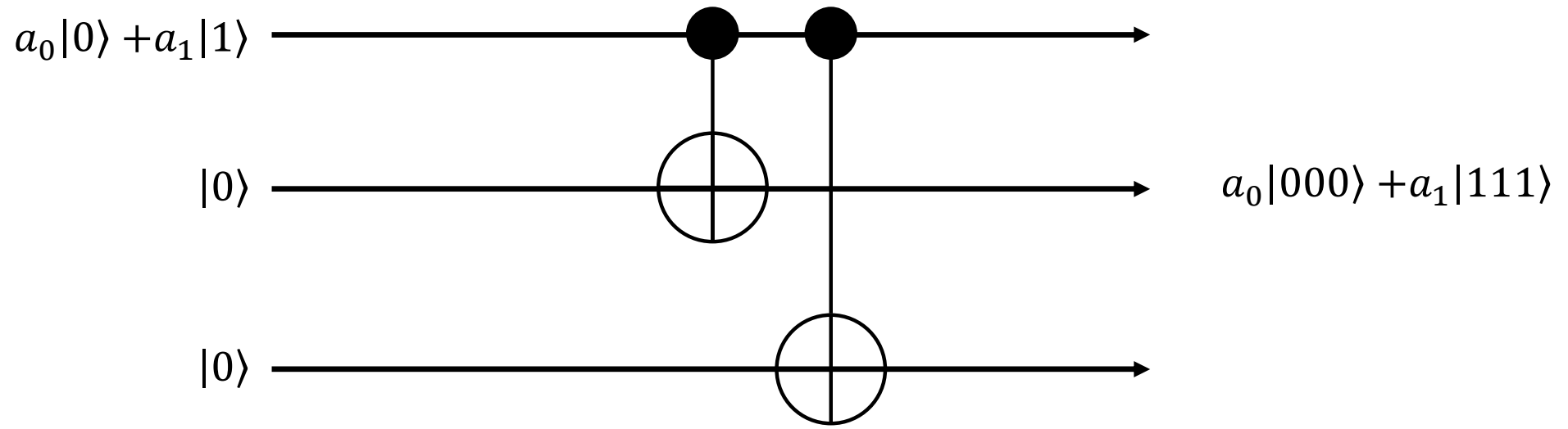
- The target input is $|y\rangle = a_0|00\rangle + a_1|11\rangle$
- The output is $|z\rangle = a_0|000\rangle + a_1|111\rangle$

Adding another CNOT



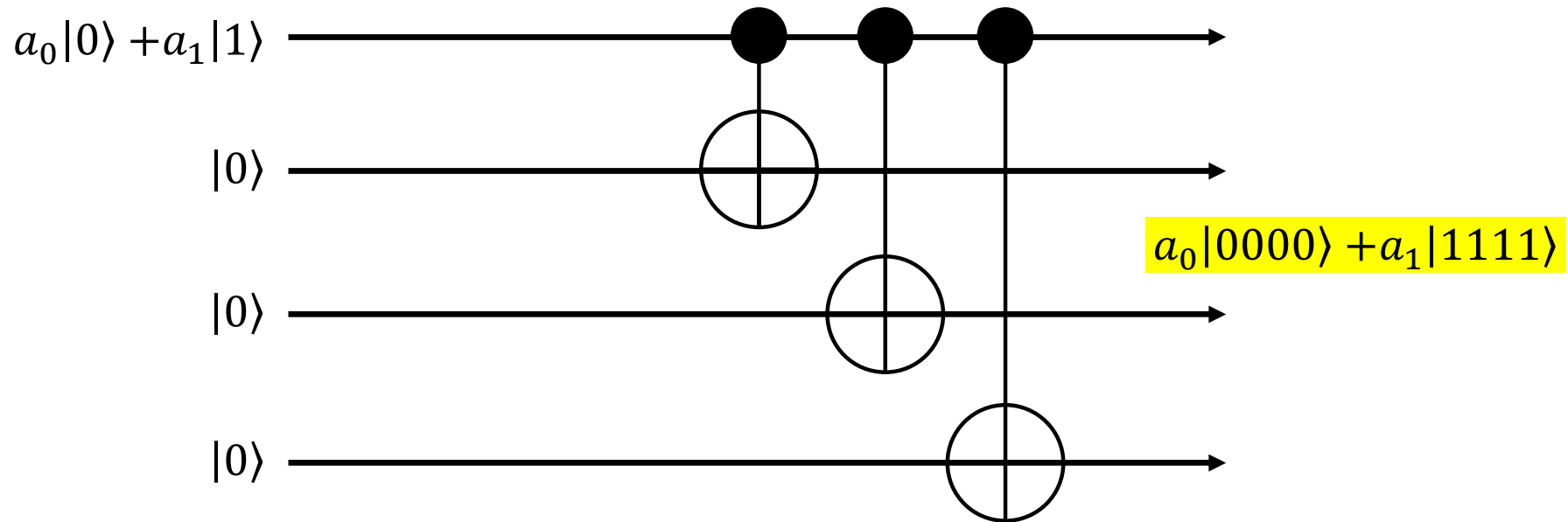
- The target input is $|y\rangle = a_0|00\rangle + a_1|11\rangle$
- The output is $|z\rangle = a_0|000\rangle + a_1|111\rangle$

Adding another CNOT



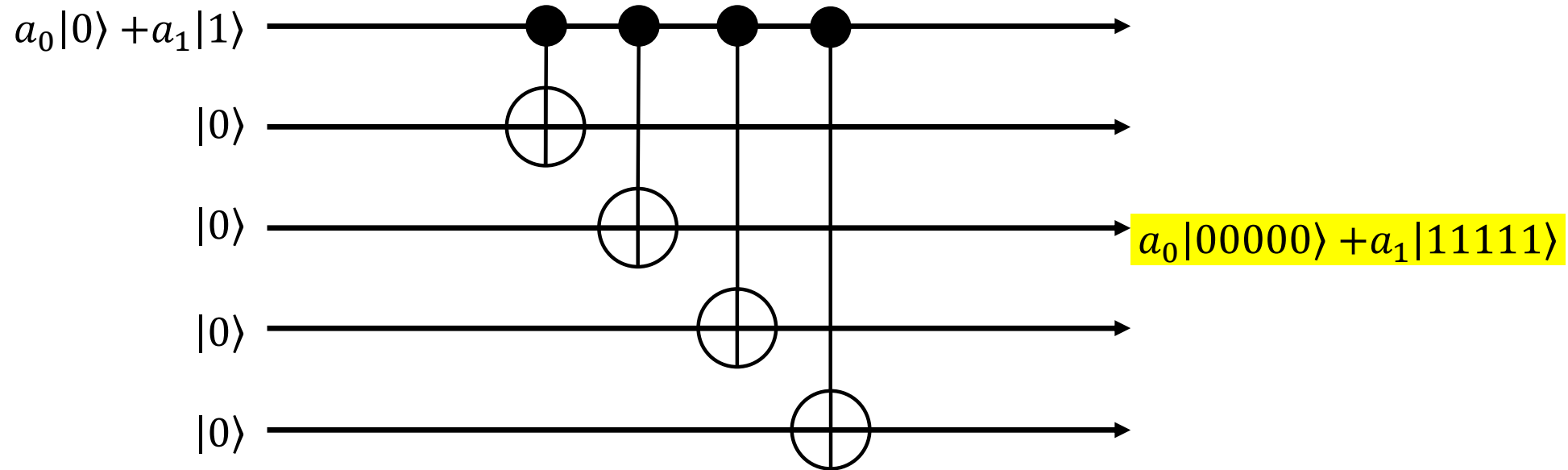
- The target input is $|y\rangle = a_0|00\rangle + a_1|11\rangle$
- The output is $|z\rangle = a_0|000\rangle + a_1|111\rangle$

Adding another CNOT



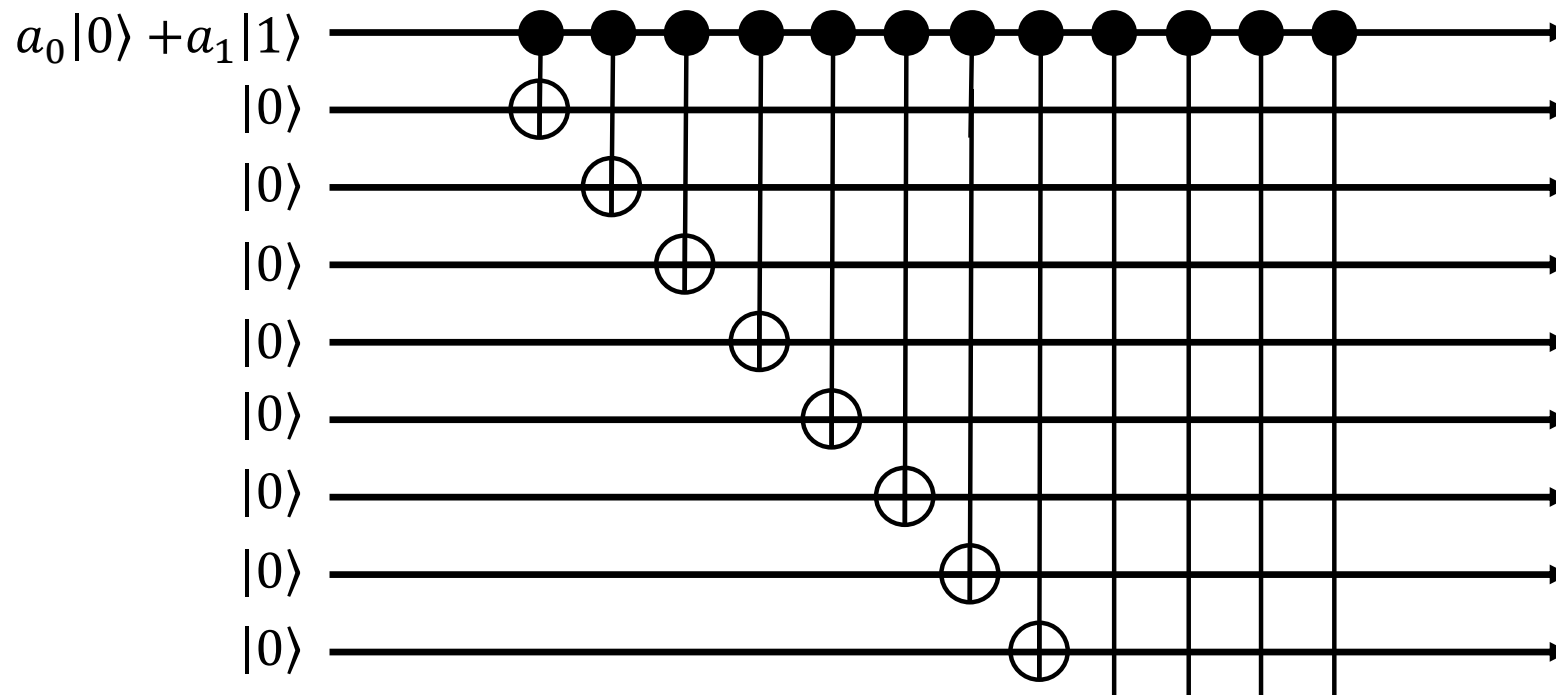
- The target input is $|z\rangle = a_0|000\rangle + a_1|111\rangle$
- The output is $|c\rangle = a_0|0000\rangle + a_1|1111\rangle$

Adding another CNOT



- The target input is $|c\rangle = a_0|00000\rangle + a_1|11111\rangle$
- The output is $|d\rangle = a_0|00000\rangle + a_1|11111\rangle$

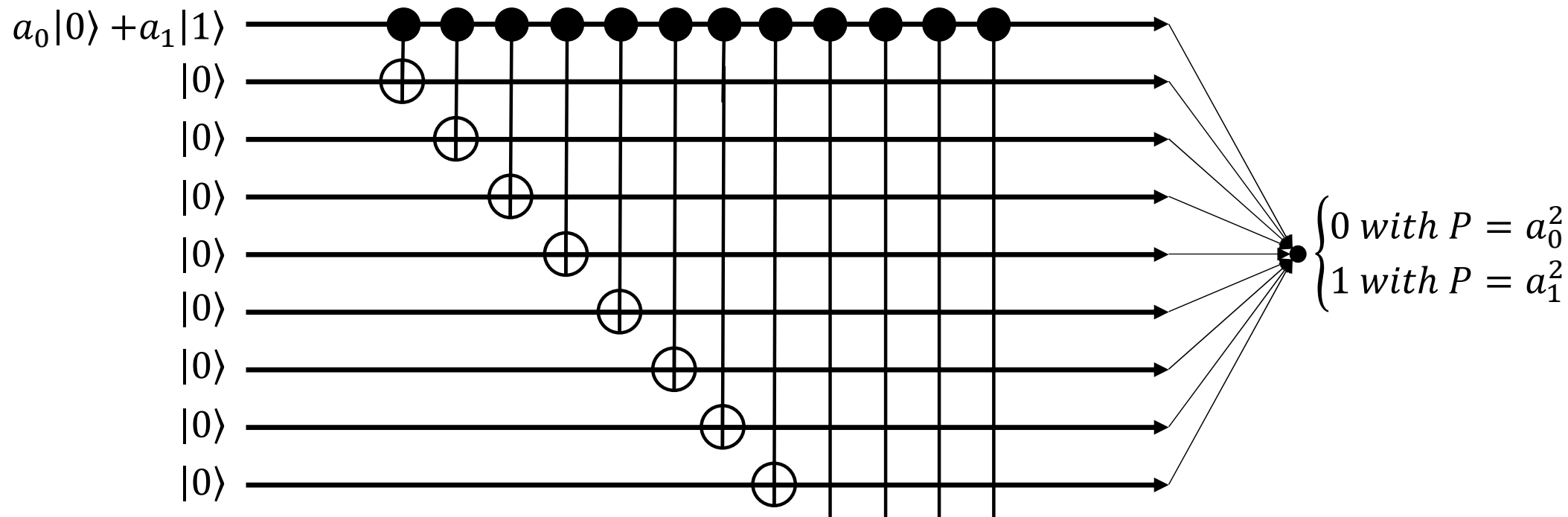
With sufficient addition we get...



- The output is

$$|o\rangle = a_0|000000 \dots 0\rangle + a_1|111111 \dots 1\rangle$$

With sufficient addition we get...



- The output is

$$|o\rangle = a_0|000000 \dots 0\rangle + a_1|111111 \dots 1\rangle$$

- Nature does not like quantum macroscopic objects
 - It will collapse this to either $|000000 \dots 0\rangle$ with probability a_0^2 or $|111111 \dots 1\rangle$ with probability a_1^2
- This gives us a measurement of 0 with $P = a_0^2$ and 1 with $P = a_1^2$

Measurement

- Previous example from Umesh Vazirani
- In general, measurement is more complex
 - But consists of composing macro quantum objects that will collapse
- More on this later
- Moving on...